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B0M33BDT – How to get/put data to Hadoop, NiFi, Kafka, Architecture patterns

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Getting Data

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Hadoop ZOO

GOVERNANCE INTEGRATION	TOOLS						SECURITY	OPERATIONS		
Data Lifecycle & Governance		Zeppelin Ambari User Views							Administration Authentication Authorization	Provisioning, Managing, & Monitoring
Falcon	DATA ACCESS							Auditing Data Protection		
Atlas Data workflow	Batch Map Reduce	Script Pig	Sql Hive	NoSql HBase Accumulo	Stream Storm	Search Solr	In-Mem Spark	Others HAWQ Partners	Ranger	Ambari Cloudbreak ZooKeeper
Sqoop		Tez	Tez	Slider				ST	Atlas HDFS Encryption	Scheduling
Flume Kafka	ume YARN: Data Operating System							Oozie		
NFS WebHDFS	HDFS Hadoop Distributed File System DATA MANAGEMENT									

How to put data to hadoop

- > Webservice webHDFS
- > SFTP & hdfs command
- > NFS gateway
- > Sqoop for databases jdbc
 - Full, increcemtal load
 - configuration via metadata
- > Flume (messages) deprecated
- > Flink streaming



Commercial solutions

- > Commercial solutions
- > Informatica, Talend, Oracle, ...
- > Typically as a bundle to something else

- > Spark and JDBC?
 - Yes, it works as well
 - But
 - Performance issues
 - Error handling



Apache Nifi

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NIFI (Niagara Files, Hortonworks Data Flow (HDF))

- NiFi is a processing engine designed to manage a continuous flow of information as a series of events in an ecosystem
- > Visual creation and management of directed graphs of processors
- > Highly concurrent model without a developer having to worry about the typical complexities of concurrency
- > Natural error handling
- > Cohesive and loosely coupled components which can then be reused



Naming Convention

- **Flowfiles** Information Packet. Represents each object moving through the system and for each one. Map object (key, value pair)
- Processors Perform the work data routing, transformation, or mediation between systems.
- Connections Connections provide the actual linkage between processors.
 These act as queues and allow various processes to interact at differing rates.
- Flow Controller Scheduler. The Flow Controller maintains the knowledge of how processes connect and manages the threads and allocations thereof which all processes use



Hello World - Example

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https://nifi.apache.org/docs.html

Architecture

- FlowFile Repository where NiFi keeps track of the state of what it knows about a given FlowFile that is presently active in the flow.
- Content Repository where the actual content bytes of a given FlowFile live
- **Provenance Repository** where all provenance event data is stored.



Cluster

-	OS,	/Host			●、	
6		0	S/Host			
<	x	₽ .	- 0	S/Host		
L		Q ₀	₿.	- 0	S/Host	
Г	V		Q0		CS/Host	ZooKeeper Server
Loc	al	V		00	G JVM Server	
	1	.ocal			C Flow Controller	
	2	Local Processo			Processor 1 Extension N	Cluster Coordinator
				Local	FlowFile Content Provenance Repository	🜲 Primary Node
					Local Storage	ZooKeeper Client

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Minifi

- A complementary data collection approach that supplements the core tenets of NiFi
- > Small size and low resource consumption
 - binary (3.2MB)
 - Original Java agent (50MB)
- Integration with NiFi for follow-on dataflow management



Apache Kafka

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Motivation for message broker

- It is a vacation day, your flight is scheduled for 14:00
- > Is my flight on time?



Getting rid of peer-to-peer data transfer



Apache Kafka



- Started as message broker at Linkedin, now it's a data processing ecosystem
- > Key characteristics
 - High-throughput
 - Distributed
 - Scalable



Kafka 101 – naming convention

- > Consumer
 - An application that is reading data from Kafka
- > Consumer Group
 - A group of an application processes that read data from Kafka
- > Producer
 - An application that is writing data to Kafka
- > Broker
 - Kafka process (single server) that is receiving data from producers, storing data on disk and provide them to consumers

Kafka 101 – naming convention

- > Topic
 - Named "message queue"
- > Partition
 - Topics are broken down into partitions
- > Offset
 - The position of a last committed message of a consumer in a topic/partition

Kafka cluster

- > Broker a single Kafka server
- > Kafka cluster collection of brokers that work together



Kafka basics – Topic and partitioning

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> Topic

- Named container for similar events
 - Usually there are many topics in a system
 - Data in one topic can be duplicated with data in another topic
- Durable logs of events
 - New message always on the end
 - Can be read by seaking arbitrary offset
 - Are immutable once something has happened, it is exceedingly difficult to make it un-happen
 - Are durable stored on filesystem

Kafka basics – Topic and partitioning

- > Partitioning
 - Breaked single topic log to multiple logs that can live independently
 - Are spread out across a cluster



Kafka basics - Offset

- > The consumer offset is a way of tracking the sequential order in which messages are received by Kafka topics.
- > Consumer offsets are persisted



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Message structure

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- > Key
- > Value
- > TopicName
- Key and Value can be any type (as long as we can serialize/deserialize them to bytes)
 - We can easily send json records as Values
 - Another very popular format to serialize data is Avro, but it requires Schema Registry

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Message ordering in Kafka

- > Message order
 - Kafka guarantees to maintain a message order **per partition**
 - That doesn't guarantee an order of messages per topic
- Messages with the same key are guaranteed to be send to the same partition
 - By default a hash partitioned is used
 - Messages without a key will be uniformly distributed between partitions

Kafka basic - Producer

- > Client application
- > Write data to appropriate kafka topic and broker
 - Serialize data
 - Define partition
 - Compress data
 - Handle errors

Acks	Throughput	Latency	Durability
0	High	Low	No Guarantee. The producer
			does not wait for
			acknowledgment from the
			server.
1	Medium	Medium	Leader writes the record to its
			local log, and responds without
			awaiting full acknowledgment
			from all followers.
-1	Low	High	Leader waits for the full set of in-
			sync replicas (ISRs) to
			acknowledge the record. This
			guarantees that the record is not
			lost as long as at least one IRS
			is active.



Kafka basic - Producing data



Kafka basic - Consumer

- > Client application
- > Read data from Kafka topic
- > Scallable organized to consumer groups
- > Keep up to date metadata (offset)



Kafka basics – Consumer Group









Kafka basics – Consumer Groups



Message delivery semantics

- > Exactly once
 - Every message is delivered only once
 - We need to keep track which messages were delivered and processed
- > At least once
 - A message might be delivered more than once
 - Might be OK given pragmatic consideration
- > At most once
 - The system will never try to deliver a message again once it was sent
 - Good option for non-critical data, that quickly become irrelevant

At least once example

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>

>

>

```
def consume_loop(consumer, topics):
   try:
        consumer.subscribe(topics)
       msg_count = 0
       while running:
            msg = consumer.poll(timeout=1.0)
            msg_process(msg)
            msg_count += 1
            consumer.commit(async=False)
   finally:
       # Close down consumer to commit final offsets.
        consumer.close()
```

At most once example

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> def consume_loop(consumer, topics):

try:

>

```
consumer.subscribe(topics)
```

```
> while running:
```

```
msg = consumer.poll(timeout=1.0)
```

consumer.commit(async=False)

```
msg_process(msg)
```

```
finally:
```

```
# Close down consumer to commit final offsets.
```

```
consumer.close()
```

Exactly once – Kafka support

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- > Idempotence: Exactly-once in order semantics per partition
 - Safeguard against duplicates in retry logic, that might be caused by broker or producer failure (hash
 - The message will be written to the Kafka topic once
 - Enable enable.idempotence=true in producer configuration
- > Transactions: Atomic writes across multiple partitions
 - New transation API atomic writes across multiple partitions
 - Consumer side configuration
 - isolation.level (read_committed, read_uncommitted)
- > Kafka streams
 - processing.guarantee=exactly_once

Log compaction

- When the data is no longer needed (after retention period) the default action is to delete the message
- Kafka has special retention policy called "compaction" in case we want to store most recent message for each key



Kafka ecosystem

Kafka Connect

- Common framework for building connectors to integrate various data stores with Kafka
- > Allows both getting data into and from Kafka
- If you find yourself trying to get data into or from Kafka, there is probably a connector for that
- > The ingest process can be speeded up by running multiple connector instances in distributed fashion
- > Aside from plain piping the data from system a to Kafka, also supports simple transformations of data records in transition

Kafka Connect- Example

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{
 "name": "gcs-sink-01",
 "config": {
 "connector.class": "io.confluent.connect.gcs.GcsSinkConnector",
 "tasks.max": "1",
 "topics": "gcs_topic",
 "gcs.bucket.name": "<my-gcs-bucket>",
 "storage.class": "io.confluent.connect.gcs.storage.GcsStorage",
 "format.class": "io.confluent.connect.gcs.format.avro.AvroFormat",
 "partitioner.class": "io.confluent.connect.storage.partitioner.DefaultPartitioner",
 "value.converter": "io.confluent.connect.avro.AvroConverter",
 "value.converter": NoNE",
 "schema.compatibility": "NONE",
 "confluent.topic.bootstrap.servers": "localhost:9092",
 "errors.tolerance": "all",
 "errors.deadletterqueue.topic.name": "dlq-gcs-sink-01"
 }

Kafka Streams

- Library build on-top of Kafka Producer/Consumer API for real-time stream processing
- > It's best suited for reading data from Kafka topic, doing some work and then writing data to another Kafka topic
- > Application instance is a JVM process
- > The parallelism of a Kafka Streams application is primarily determined by how many partitions the input topics have



Kafka Streams naming convention

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> Source processor

- produces an input stream to its topology from one or multiple Kafka topics by consuming records from these topics.
- > Sink processor
 - sends any received records from its up-stream processors to a specified Kafka topic.
- > Stream processor
 - represents a processing step in a topology it is
 Standard operations are map or filter, joins, and
- > Stream task
 - smallest unit of work within a Kafka Streams
 application instance. The number of tasks is determined
 by an application's source topic with the highest number
 of partitions



Kafka Streams Architecture



ksqIDB



- > DB-like abstraction on top of Kafka Streams
- Provides table-like interface over Kafka topic (using extended SQL syntax)
- > Main components
 - KSQL Server processes SQL statements and queries
 - KSQL CLI CLI program to i ksqIDB architecture and components



Kafka use-cases

- > IoT and any message-oriented application
 - Sensor data
 - Financial transactions
 - Stock market
 - Logs
- > Asynchronous application communication
- > Publish subscribe messaging
- > Data integration (Kappa & Lambda architecture)

Architectures

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Lambda

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- > From Apache Storm
- > Nathan Marz, 2011

Yahoo, Netflix

> http://nathanmarz.com/blog/how-to-beat-the-cap-theorem.html



Lambda

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> 4 layers



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Lambda

> Mapping to the technologies



Kappa



- > 2014 Jay Kreps Linkedin
- > https://www.oreilly.com/ideas/questioning-the-lambda-architecture



Kappa

- > 3 layers batch layer removed
- > Long retention can be used
- > How to work with state microbatches?



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Kappa

> Mapping to technologies





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