

# Assignment #1-3 Consultations

Exercises on heuristic computation

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PUI Tutorial  
Week 8

# Assignment #1-3 Consultations

When the assignment is due  
Friday but today is  
Wednesday.



# Relaxation heuristics

$$F = \{a, b, c, d, e, f\}$$

$$s_{init} = \{a, b\}$$

$$s_{goal} = \{d, f\}$$

	pre	add	del	c	
$O =$	$o_1$	a,b	b	a	1
	$o_2$	b,c	a	b	1
	$o_3$	c	e	$\emptyset$	2
	$o_4$	c,e	d	e	2
	$o_5$	a,c,d	f	a	1
	$o_6$	a,b	c	$\emptyset$	1

Compute

- $h^{add}(s_{init})$
- $h^{max}(s_{init})$

**Algorithm 1:** Algorithm for computing  $h^{\max}(s)$ .**Input:**  $\Pi = \langle \mathcal{F}, \mathcal{O}, s_{init}, s_{goal}, c \rangle$ , state  $s$ **Output:**  $h^{\max}(s)$ 

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1 for each  $f \in s$  do  $\Delta_1(s, f) \leftarrow 0$ ;
2 for each  $f \in \mathcal{F} \setminus s$  do  $\Delta_1(s, f) \leftarrow \infty$ ;
3 for each  $o \in \mathcal{O}, pre(o) = \emptyset$  do
4   | for each  $f \in add(o)$  do  $\Delta_1(s, f) \leftarrow \min\{\Delta_1(s, f), c(o)\}$ ;
5 end
6 for each  $o \in \mathcal{O}$  do  $U(o) \leftarrow |pre(o)|$ ;
7  $C \leftarrow \emptyset$ ;
8 while  $s_{goal} \not\subseteq C$  do
9   |  $k \leftarrow \arg \min_{f \in \mathcal{F} \setminus C} \Delta_1(s, f)$ ;
10  |  $C \leftarrow C \cup \{k\}$ ;
11  | for each  $o \in \mathcal{O}, k \in pre(o)$  do
12  |   |  $U(o) \leftarrow U(o) - 1$ ;
13  |   | if  $U(o) = 0$  then
14  |   |   | for each  $f \in add(o)$  do
15  |   |   |   |  $\Delta_1(s, f) \leftarrow \min\{\Delta_1(s, f), c(o) + \Delta_1(s, k)\}$ ;
16  |   |   |   end
17  |   |   end
18  |   end
19 end
20  $h^{\max}(s) = \max_{f \in s_{goal}} \Delta_1(s, f)$ ;

```

# Relaxation heuristics

$$F = \{a, b, c, d, e, f\}$$

$$s_{init} = \{a, f\}$$

$$s_{goal} = \{d\}$$

$$O =$$

	pre	add	del	c
$o_1$	a,f	b	f	1
$o_2$	c	a,d	c	1
$o_3$	a,b	c,e	a	1
$o_4$	c,b	d	c,b	2
$o_5$	a	e,f	$\emptyset$	3

Compute

- $h^{LM-Cut}(s_{init})$
- $h^{FF}(s_{init})$

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**Algorithm 2:** Algorithm for computing  $h^{lm-cut}(s)$ .

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**Input:**  $\Pi = \langle \mathcal{F}, \mathcal{O}, s_{init}, s_{goal}, c \rangle$ , state  $s$

**Output:**  $h^{lm-cut}(s)$

- 1 **if**  $h^{\max}(\Pi, s_{init}) = \infty$  **then**
  - 2 |  $h^{lm-cut}(s) \leftarrow \infty$  and terminate;
  - 3 **end**
  - 4  $h^{lm-cut}(s) \leftarrow 0$ ;
  - 5  $\Pi_1 = \langle \mathcal{F}' = \mathcal{F} \cup \{I, G\}, \mathcal{O}' = \mathcal{O} \cup \{o_{init}, o_{goal}\}, s'_{init} = \{I\}, s'_{goal} = \{G\}, c_1 \rangle$ , where  
 $\text{pre}(o_{init}) = \{I\}$ ,  $\text{add}(o_{init}) = s$ ,  $\text{del}(o_{init}) = \emptyset$ ,  $\text{pre}(o_{goal}) = s_{goal}$ ,  $\text{add}(o_{goal}) = \{G\}$ ,  
 $\text{del}(o_{goal}) = \emptyset$ ,  $c_1(o_{init}) = 0$ ,  $c_1(o_{goal}) = 0$ , and  $c_1(o) = c(o)$  for all  $o \in \mathcal{O}$ ;
  - 6  $i \leftarrow 1$ ;
  - 7 **while**  $h^{\max}(\Pi_i, s'_{init}) \neq 0$  **do**
  - 8 | Construct a justification graph  $G_i$  from  $\Pi_i$ ;
  - 9 | Construct an s-t-cut  $\mathcal{C}_i(G_i, n_I, n_G) = (N_i^0, N_i^* \cup N_i^b)$ ;
  - 10 | Create a landmark  $L_i$  as a set of labels of edges that cross the cut  $\mathcal{C}_i$ , i.e., they  
lead from  $N_i^0$  to  $N_i^*$ ;
  - 11 |  $m_i \leftarrow \min_{o \in L_i} c_i(o)$ ;
  - 12 |  $h^{lm-cut}(s) \leftarrow h^{lm-cut}(s) + m_i$ ;
  - 13 | Set  $\Pi_{i+1} = \langle \mathcal{F}', \mathcal{O}', s'_{init}, s'_{goal}, c_{i+1} \rangle$ , where  $c_{i+1}(o) = c_i(o) - m_i$  if  $o \in L_i$ , and  
 $c_{i+1}(o) = c_i(o)$  otherwise;
  - 14 |  $i \leftarrow i + 1$ ;
  - 15 **end**
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## Overall algorithm:

- Create reachability graph
- Mark the final G node
- Apply rules layers by layer until every marked node is justified

## Justified node definitions:

- Action node is justified if all precondition fact nodes are marked
- Fact node is justified if at least one predecessor node is marked
  - Starting with marked goal node, apply the following rules **layer by layer** until **all marked nodes are justified**
    - 1) Mark all immediate predecessors of a marked unjustified action node
    - 2) Mark the immediate predecessor of a marked unjustified atom node with only one immediate predecessor
    - 3) Mark an immediate predecessor of a marked unjustified atom node connected via an idle arc (to the same atom in the previous layer)
    - 4) Mark any immediate predecessor of a marked unjustified atom node

$$V = \{v_1, v_2, v_3\}$$

$$D_{v_1} = \{A, B\}, D_{v_2} = \{C, D\}, D_{v_3} = \{E, F\}$$

$$S_{init} = \{v_1 = A, v_2 = C, v_3 = E\}$$

$$S_{goal} = \{v_3 = F\}$$

	pre	eff	c
$O =$	$v_1 = A$	$v_1 = B$	3
$o_2$	$v_1 = B, v_2 = C$	$v_1 = A, v_2 = D$	1
$o_3$	$v_1 = A, v_2 = D$	$v_3 = F$	1
$o_4$	$v_2 = D, v_3 = E$	$v_2 = C$	2

Compute

- $h^{flow}$
- $h^{pot}$
- $h^{M\&S}$



## LP formulation

$$\text{minimize } \sum_{o \in O} c(o)x_o$$

$$\text{subject to } LB_{V,v} \leq \sum_{o \in \text{prod}(\langle V,v \rangle)} x_o - \sum_{o \in \text{cons}(\langle V,v \rangle)} x_o, \forall V \in \mathbf{V}, \forall v \in D_V$$

$$\text{where } LB_{V,v} = \begin{cases} 0 & \text{if } V \in \text{vars}(s_{goal}) \text{ and } s_{goal}[V] = v \text{ and } s[V] = v, \\ 1 & \text{if } V \in \text{vars}(s_{goal}) \text{ and } s_{goal}[V] = v \text{ and } s[V] \neq v, \\ -1 & \text{if } (V \notin \text{vars}(s_{goal}) \text{ or } s_{goal}[V] \neq v) \text{ and } s[V] = v, \\ 0 & \text{if } (V \notin \text{vars}(s_{goal}) \text{ or } s_{goal}[V] \neq v) \text{ and } s[V] \neq v, \end{cases}$$

## LP formulation

$$\text{maximize } \sum_{V \in \mathbf{V}} P_{V, s_{init}[V]}$$

$$\text{subject to } P_{V, v} \leq M_V, \forall V \in \mathbf{V}, \forall v \in D_V$$

$$\sum_{V \in \mathbf{V}} \text{maxpot}(V, s_{goal}) \leq 0$$

$$\sum_{V \in \text{vars}(\text{eff}(o))} (\text{maxpot}(V, \text{pre}(o)) - P_{V, \text{eff}(o)[V]}) \leq c(o), \forall o \in O$$

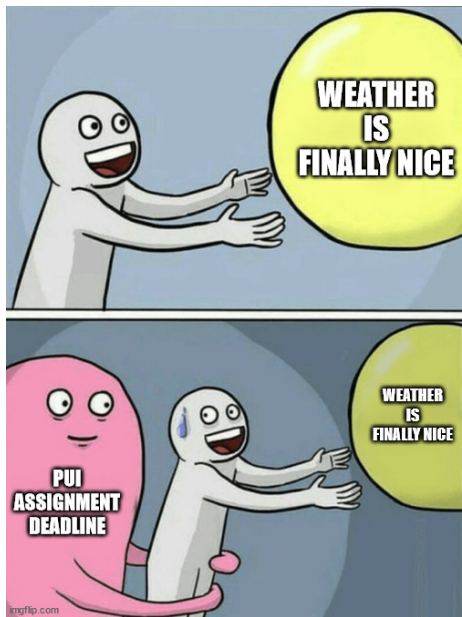
$$\text{where } \text{maxpot}(V, p) = \begin{cases} P_{V, p[V]} & \text{if } V \in \text{vars}(p), \\ M_V & \text{otherwise.} \end{cases}$$

The **value of  $h^{pot}$  heuristic** for the state  $s$  is

$$h^{pot}(s) = \begin{cases} \sum_{V \in \mathbf{V}} P_{V, s[V]} & \text{if the solution is } \mathbf{feasible} \\ \infty & \text{if the solution is } \mathbf{not\ feasible} \end{cases}$$

- ① Create atomic projections (one per variable)
- ② Merge two arbitrary transition systems (synchronized product)
- ③ Shrink the new transition graph (merge states together to create smaller abstraction)
- ④ Repeat 2 and 3 until you're left with one abstraction in which you can find the solution

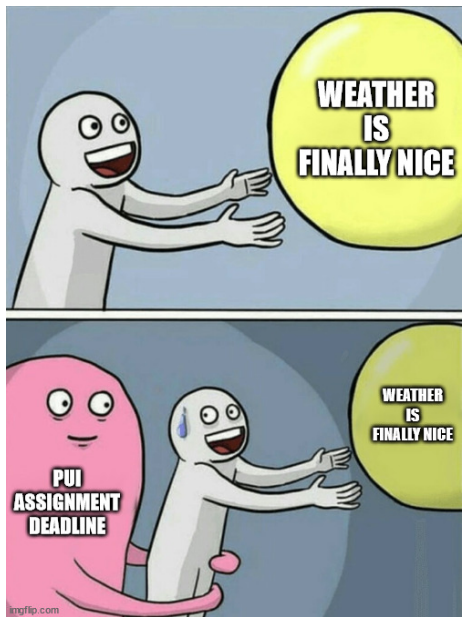
# The End



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# The End



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