

Polynomials

We will consider polynomials in n unknowns x_1, x_2, \dots, x_n with rational coefficients a_1, a_2, \dots, a_n

Polynomials are linear combinations of a finite number of monomials

Monomials : $x_1^{\alpha_1} x_2^{\alpha_2} \dots x_n^{\alpha_n}$, where α_i are non-negative integers
 $\alpha_i \in \mathbb{Z}_{\geq 0}$ - exponents

To simplify the notation we will write

$$x^\alpha = x_1^{\alpha_1} \dots x_n^{\alpha_n}$$

for n -tuple of exponents $\alpha = (\alpha_1, \alpha_2, \dots, \alpha_n)$

n -tuple α is called multidegree of monomial x^α

e.g. for $\alpha = (2, 0, 1)$ we get $x^\alpha = x_1^2 x_2^0 x_3 = x_1^2 x_3$

We define the total degree of a non-zero monomial with exponent

$\alpha = (\alpha_1, \dots, \alpha_n)$ as $d = \alpha_1 + \alpha_2 + \dots + \alpha_n$

The total degree $\deg(f)$ of a polynomial f is the maximum of the total degrees of its monomials (zero polynomial has no degree)

With this notation polynomials with rational coefficients can be written in the form:

$$f = \sum_{\alpha} a_{\alpha} x^{\alpha} \quad a_{\alpha} \in \mathbb{Q}$$

where the sum is over a finite set of n -tuples $\alpha \in \mathbb{Z}_{\geq 0}^n$

The set of all polynomials in unknowns x_1, x_2, \dots, x_n with rational coefficients will be denoted $\mathbb{Q}[x_1, \dots, x_n]$

There is an infinite (countable) number of monomials

They can be ordered using a total ordering (linear ordering) (ordering where every two elements are comparable)

Polynomials with rational coefficients can be also understood as complex functions
We can evaluate polynomial f on point $\vec{p} \in \mathbb{C}^n$

$$f(\vec{p}) = \left(\sum_{\alpha} a_{\alpha} x^{\alpha} \right) (\vec{p}) = \sum_{\alpha} a_{\alpha} x^{\alpha} (\vec{p}) = \sum_{\alpha} a_{\alpha} \vec{p}^{\alpha} = \sum_{\alpha} a_{\alpha} p_1^{\alpha_1} p_2^{\alpha_2} \dots p_n^{\alpha_n}$$

Evaluated polynomial is a linear combination of the evaluated monomials

Division of terms

$$\alpha, \beta \in \mathbb{Z}_{\geq 0}^n$$

$$a_{\alpha}, b_{\alpha} \in \mathbb{Q} [k]$$

$$x^{\alpha}, x^{\beta} \in \mathbb{Q} [x_1, \dots, x_n] \quad \text{- monomials}$$

$$a_{\alpha} x^{\alpha} \text{ divides } b_{\beta} x^{\beta} \stackrel{\text{def}}{\equiv} \beta_i - \alpha_i \geq 0, \quad i=1, \dots, n$$

If $a_{\alpha} x^{\alpha}$ divides $b_{\beta} x^{\beta}$ then there is exactly one monomial

$$c_{\gamma} x^{\gamma} = \frac{b_{\beta}}{a_{\alpha}} x^{\beta - \alpha}$$

For monomials in one variable $\alpha, \beta \in \mathbb{Z}_{\geq 0}$

Univariate polynomials

Polynomials in a single unknown are often called univariate polynomials

\Rightarrow - α is a single number

- The total degree $\deg(f)$ of f is then called degree

- The set of all polynomials in a single unknown X over rational numbers $\mathbb{Q}[X]$ forms a ring

- Polynomials are almost as real numbers except for the division

- Polynomials can't in general be divided

- In fact polynomials behave in many aspects as whole numbers \mathbb{Z}

It is easy to introduce long polynomial division in the same way as it is used in whole numbers

Leading term : of a non-zero polynomial

$$f(x) = a_m x^m + a_{m-1} x^{m-1} + \dots + a_0 \in \mathbb{Q}[x]$$

$\neq 0$

$$\text{LT}(f) = a_m x^m \equiv \text{leading term}$$

Example: $f(x) = 2x^3 - 4x + 3 \Rightarrow \text{LT}(f) = 2x^3$

"Division theorem"

Consider polynomials $f, g \in \mathbb{Q}[x]$ $g \neq 0$

Then there are unique polynomials $q, r \in \mathbb{Q}[x]$ such that

$$f = q \cdot g + r \quad \text{with} \quad \deg(r) < \deg(g) \\ \text{or} \quad r = 0$$

q - is the quotient

r - is the remainder (of f on division by g)

One often writes $f \equiv r \pmod{g}$ ($r = f \pmod{g}$)

Example: $f = 2x^3 - 4x + 3$ $g(x) = x - 1$

$$f : g = (2x^2 + 2x - 2)(x - 1) + 1$$

Notice $\deg(f) = \deg(\text{LT}(f))$

$\text{LT}(g)$ divides $\text{LT}(f) \Leftrightarrow \deg(\text{LT}(g)) \leq \deg(\text{LT}(f)) \Leftrightarrow \deg(g) \leq \deg(f)$

Proof: of Division theorem \Rightarrow Division algorithm

Input: g, f

Output: q, r

$q := 0$

$r := f$

WHILE $r \neq 0$ AND $\text{LT}(g)$ divides $\text{LT}(r)$

{ $q := q + \frac{\text{LT}(r)}{\text{LT}(g)}$

$r := r - \frac{\text{LT}(r)}{\text{LT}(g)} \cdot g$

}

$$f = 2x^3 - 4x + 3$$

$$g = x - 1$$

$$q = 0$$

$$r = 2x^3 - 4x + 3$$

$$q = 2x^2$$

$$r = 2x^2 - 4x + 3$$

$$q = 2x^2 + 2x$$

$$r = -2x + 3$$

$$q = 2x^2 + 2x - 2$$

$$r = 1$$

$$f = (2x^2 + 2x - 2)(x - 1) + 1$$

Observe that $f = qg + r$ holds true

$$a) \quad q=0 \quad \& \quad r=f \quad \Rightarrow \quad 0 \cdot g + f = f$$

b) Let q_i, r_i be such that $f = q_i g + r_i$, then

$$q_{i+1} g + r_{i+1} = \underbrace{\left(q_i + \frac{LT(r_i)}{LT(g)} \right)}_{q_{i+1}} g + \underbrace{\left(r_i - \frac{LT(r_i)}{LT(g)} \cdot g \right)}_{r_{i+1}}$$

$$= q_i g + r_i = f$$

If the algorithm terminates, then either

$$r=0 \quad \text{or} \quad LT(g) \text{ does not divide } LT(r) \Leftrightarrow \deg(r) < \deg(g)$$

Let us show that the algorithm terminates

Assume that the algorithm does not terminate
 then $LT(g)$ divides $LT(r)$ and $r \neq 0$

Observe that for $r_{i+1} = r_i - \frac{LT(r_i)}{LT(g)} \cdot g$ holds $r_{i+1} \begin{cases} \text{either} = 0 \\ \text{or } \deg(r_{i+1}) < \deg(r_i) \end{cases}$

write $r_i = a_0 x^m + a_1 x^{m-1} + \dots + a_m$
 $g = b_0 x^l + b_1 x^{l-1} + \dots + b_l$

with $m \geq l$ ($LT(g)$ divides $LT(r_i)$)

$$r_{i+1} = r_i - \frac{LT(r_i)}{LT(g)} g = \underbrace{(a_0 x^m + \dots + a_m)}_{\text{cancel}} - \frac{a_0}{b_0} x^{m-l} \underbrace{(b_0 x^l + \dots)}_{\text{cancel}}$$

$$= (a_1 x^{m-1} + \dots) - \left(\frac{a_0}{b_0} b_1 x^{m-1} + \dots \right)$$

$$= \left(a_1 - \frac{a_0}{b_0} b_1 \right) x^{m-1} + \left(a_2 - \frac{a_0}{b_0} b_2 \right) x^{m-2} + \dots$$

and therefore

- either $r_{i+1} = 0$ (if all coefficients vanish)
- or $\deg(r_{i+1}) \leq m-1 < m = \deg(r_i)$

Monomials in one variable are easy to order by their degree, i.e

$$x^0 <_{\text{deg}} x^1 <_{\text{deg}} x^2 \dots$$

also notice $x^m <_{\text{deg}} x^n \Leftrightarrow x^m$ divides x^n

Not so simple with more variables

Consider xy^2, x^2y - neither one divides the other

$$\deg(xy^2) = 1+2 = 3 = 2+1 = \deg(x^2y)$$

total degrees does not define an ordering of monomials

by degrees \rightarrow partial ordering

to have a total ordering we have to extend this

A monomial ordering on $k[x_1, \dots, x_n]$ is any ordering relation $<$ on \mathbb{Z}_{\geq}^n satisfying

(i) $\forall \alpha, \beta : \alpha > \beta$ or $\alpha < \beta$

(ii) $\alpha > \beta$ & $\gamma \in \mathbb{Z}_{\geq 0}^n \Rightarrow \alpha + \gamma > \beta + \gamma$ ($\deg(m_1) < \deg(m_2)$
 $\Rightarrow \deg(m_1 \cdot m) < \deg(m_2 \cdot m)$)

(iii) $\forall \alpha : \alpha > 0$

We write $x^\alpha > x^\beta \stackrel{\text{def.}}{\equiv} \alpha > \beta$

Lexicographic ordering

$\alpha = (\alpha_1, \alpha_2, \dots, \alpha_n)$, $\beta = (\beta_1, \beta_2, \dots, \beta_n) \in \mathbb{Z}_{\geq 0}^n$

$\alpha >_{\text{lex}} \beta$ if the left-most non-zero element of $\alpha - \beta$ is positive or $\alpha - \beta = 0$

Example : $(1, 2, 0) >_{\text{lex}} (0, 3, 4) \quad \equiv \quad (1, -1, -4)$
 $(3, 2, 4) >_{\text{lex}} (3, 2, 1) \quad \equiv \quad (0, 0, 3)$

Lex orders monomials as words in dictionary

An important parameter of $<_{\text{lex}}$ ordering is the order of unknowns (order of letters)

e.g. $xy^2z = x.y.y.z <_{\text{lex}} x.y.z.z = xyz^2$ if $x < y < z$

$$xyz^2 = xyzz <_{\text{lex}} xyzy = xy^2z \quad \text{if } z < y < x$$

There are $n!$ possible $<_{\text{lex}}$ orderings when dealing with n unknowns

Proof: The lex ordering on \mathbb{Z}_{\geq}^n is a monomial ordering

1. $<_{\text{lex}}$ is an ordering ($\alpha > \alpha$; $\alpha > \beta$ & $\beta > \gamma \Rightarrow \alpha > \gamma$
 $\alpha > \beta$ & $\beta > \alpha \Rightarrow \alpha = \beta$)

a) $\alpha - \alpha = 0 \Rightarrow \alpha >_{\text{lex}} \alpha$

$$b) \quad \alpha >_{\text{lex}} \beta, \quad \beta >_{\text{lex}} \gamma$$

$$\Rightarrow \exists i, j \in \mathbb{Z}_{\geq 0}^n \text{ such that}$$

$$(\alpha - \beta)_h = 0 \text{ and } (\beta - \gamma)_m = 0 \text{ for } h < i, m < j \text{ \& } (\alpha - \beta)_i > 0 \text{ \& } (\beta - \gamma)_j > 0$$

$$(\alpha - \gamma)_h = 0 \quad h = 1, \dots, (\min(i, j)) - 1 \quad \alpha_h = \beta_h = \gamma_h$$

$$(\alpha - \gamma)_{\min(i, j)} > 0 \quad \begin{cases} \min(i, j) = i \\ \min(i, j) = j \end{cases} \quad \begin{matrix} \alpha_i > \beta_i = \gamma_i \\ \alpha_j = \beta_j > \gamma_j \end{matrix}$$

$$\Rightarrow \alpha >_{\text{lex}} \gamma$$

$$c) \quad \alpha >_{\text{lex}} \beta \text{ \& } \beta >_{\text{lex}} \alpha \Rightarrow \text{either } \alpha - \beta = 0$$

$$\text{or } \exists i \in \mathbb{Z}_{\geq 0} \quad (\alpha - \beta)_i > 0 \Rightarrow \alpha - \beta = 0 \\ \text{\& } (\beta - \alpha)_i > 0$$

The lex ordering is a monomial ordering

i) $\forall \alpha, \beta \quad \alpha >_{\text{lex}} \beta \quad \text{or} \quad \beta >_{\text{lex}} \alpha :$

$c = \alpha - \beta = 0 \Rightarrow \alpha = \beta$ or there is the first non-zero element c_i . If $c_i > 0$ then $\alpha >_{\text{lex}} \beta$ otherwise $\beta >_{\text{lex}} \alpha$

ii) $\alpha >_{\text{lex}} \beta \quad \& \quad \gamma \in \mathbb{R}_{\geq}^n \Rightarrow \alpha + \gamma >_{\text{lex}} \beta + \gamma$

$$(\alpha + \gamma) - (\beta + \gamma) = \alpha - \beta$$

iii) $\forall \alpha : \alpha >_{\text{lex}} 0$

$$(\alpha - 0)_i \geq 0$$

Graded reverse Lex monomial ordering (Grevlex) $<_{\text{grevlex}}$

is an extension of the partial ordering by the total degree to a total monomial ordering

Monomial $x^\alpha <_{\text{grevlex}} x^\beta$ (as well as $\alpha <_{\text{grevlex}} \beta$ for exponents) when

either $\deg(\alpha) < \deg(\beta)$ or $\deg(\alpha) = \deg(\beta)$ and the last non-zero
element of $\beta - \alpha$ is negative

Polynomial division with more divisors in more variables

Division by more than one polynomial

$$f = 3x^4 - x^2 + 2x, \quad f_1 = x-1, \quad f_2 = x^2+1$$

$$\begin{aligned} f &= 0 \cdot (x-1) + 0(x^2+1) + 3x^4 - x^2 + 2x && + 0 \\ &= 3x^3(x-1) + 0(x^2+1) + 3x^3 - x^2 + 2x && + 0 \\ &= (3x^3 + 3x^2)(x-1) + 0(x^2+1) + 2x^2 + 2x && + 0 \\ &= (3x^3 + 3x^2 + 2x)(x-1) + 0(x^2+1) + 4x && + 0 \\ &= (3x^3 + 3x^2 + 2x + 4)(x-1) + 0(x^2+1) && + 4 \end{aligned}$$

$$\begin{aligned} &= 3x(x^2+1) + 0(x-1) - x^2 - x && + 0 \\ &= (3x-1)(x^2+1) + 0(x-1) - x + 1 && + 0 \\ &= (3x-1)(x^2+1) - 1(x-1) && + 0 \end{aligned}$$

We see that $f: (f_1, f_2) \neq f: (f_2, f_1) \Rightarrow f: \{f_1, f_2\}$ is not well defined

"Division theorem" for more than one divisor in $k[x_1, \dots, x_n]$

Let $>$ be a monomial order on $\mathbb{Z}_{\geq 0}^n$ and $F = (f_1, \dots, f_s)$ an ordered s -tuple of $f_i \in k[x_1, \dots, x_n]$. Then every $f \in k[x_1, \dots, x_n]$ can be written as:

$$f = a_1 f_1 + \dots + a_s f_s + r$$

$a_i, r \in k[x_1, \dots, x_n]$ and

either $r = 0$

or none of the monomials of r is divisible by any of $LT(f_1), \dots, LT(f_s)$

Furthermore $a_i f_i \neq 0 \Rightarrow \text{multidegree}(f) \geq \text{multidegree}(a_i f_i)$

$r \equiv$ remainder of f on division by F $r = \bar{f}^F$

"Division algorithm" for more than one divisor in $k[x_1, \dots, x_n]$

Input : $F = (f_1, \dots, f_s), f \in F$

Output : $a_1, \dots, a_s, r \in F$

$a_1 := a_2 := \dots := a_s := r := 0, p := f$

WHILE $p \neq 0$ DO

{ $i := 1$

 divisionoccured := FALSE

 WHILE $i \leq s$ AND divisionoccured = FALSE DO

 { IF $LT(f_i)$ divides $LT(p)$ THEN

 { $a_i := a_i + \frac{LT(p)}{LT(f_i)}$

$p := p - \frac{LT(p)}{LT(f_i)} \cdot f_i$

 divisionoccured := TRUE }

 ELSE { $i := i + 1$ } }

 IF divisionoccured = FALSE THEN

 { $r := r + LT(p)$

$p := p - LT(p)$ }

}

[Proof as for 1 variable
degree \rightarrow multidegree
 $r \rightarrow p$]

Example

$$x >_{\text{lex}} y \quad f = xy^2 + x + 1, \quad f_1 = xy + 1, \quad f_2 = y + 1$$

multidegrees

$\downarrow \quad \downarrow \quad \downarrow$
 $(1,2) \quad (1,0) \quad (0,0)$

$\downarrow \quad \downarrow$
 $(1,1) \quad (0,0)$

$\downarrow \quad \downarrow$
 $(0,1) \quad (0,0)$

$$f = y \cdot (xy + 1) + \underbrace{x - y + 1}_{r} = y(xy + 1) - 1(y + 1) + x + 2$$

$\downarrow \quad \downarrow \quad \downarrow$
 $(1,0) \quad (0,1) \quad (0,0)$

\downarrow
 a_1

\downarrow
 a_2

$$f = \underbrace{0}_{a_1} \cdot f_1 + \underbrace{0}_{a_2} \cdot f_2 + \overbrace{xy^2 + x + 1}^p + \underbrace{0}_r$$

$$= y \cdot f_1 + 0 \cdot f_2 + x - y + 1 + 0$$

$$= y \cdot f_1 + 0 \cdot f_2 + -y + 1 + x$$

$$= y \cdot f_1 - 1 \cdot f_2 + 2 + x$$

$$= y \cdot f_1 - 1 \cdot f_2 + x + 2$$

The order of monomials in F matters!

$$f = xy^2 - x, \quad f_1 = xy + 1, \quad f_2 = y^2 - 1 \quad >_{\text{lex}}, \quad x >_{\text{lex}} y$$

a) $f: (f_1, f_2)$

$$xy^2 - x = \underbrace{y}_{a_1} (xy + 1) + \underbrace{0}_{a_2} \cdot (y^2 - 1) + \underbrace{(-x - y)}_r$$

b) $f: (f_2, f_1)$

$$xy^2 - x = \underbrace{x}_{a_1} (y^2 - 1) + \underbrace{0}_{a_2} \cdot (xy + 1) + \underbrace{0}_r$$