Roboter Navigation Temporal Task-Motion Planning Stefan Edelkamp

Driving Research Questions

How can we improve motion planning for complex systems?

- How can we develop motion planners that are generally applicable?
- How can we achieve planning efficiency even with nonlinear dynamics?
- How far back can we push the "curse of dimensionality"?
- Is there Pareto optimality between efficiency and solution quality?
- What formal guarantees can we provide?

Framework

- Sampling-based motion planning
 - \Rightarrow generality: dynamics as *black-box function* $s_{new} \leftarrow MOTION(s, u, dt)$
 - ⇒ continuous state/control spaces: *probabilistic sampling to make it feasible*
 - \Rightarrow high-dimensionality: *search to find solution*

coupled with discrete abstractions

- \Rightarrow provide simplified planning layer
- \Rightarrow guide search in the continuous state/control spaces

and motion controllers

- \Rightarrow open up the black-box MOTION function
- \Rightarrow facilitate search expansion

Formal guarantees

 \Rightarrow probabilistic completeness

For a Start: The Physical TSP



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Traveling Salesman Tours (TSP)

- Given a map, compute a minimum-cost round trip visiting certain cities
- Shortest paths graph reduction: precompute all-pairs-shortest-paths with
- Dijkstra's algorithm (be smart: employ radix heaps)
- Traditional: Model problem as an IP and call solver (CPLEX, IPSolve, . . .)
- Neighborhood search (xOPT: SA; GA; AA; PSO; LNS,...)
- . . . New in the arena: Monte-Carlo Search





Additional Constraints

Time Windows, Capacities, Premium Services, Pickup and Deliveries TSP+TW: Restricted time intervals / service times C+TSP: Limited vehicle load

TSP+PD: Pickup and deliveries (PDP)



TSP*: Premium service - same-day delivery preferred

VRP: Vehicle routing - several vehicles



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Inspection Problem



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Inspection Tour



Grassfiring



Grassfiring





Generating Inspection Waypoints (1)

Input: \mathcal{I} : bitmap image; α : desired inspection quality, $0 < \alpha \leq 1$ Output: a set of inspection points

1:
$$h \leftarrow \operatorname{height}(\mathcal{I})$$
; $w \leftarrow \operatorname{width}(\mathcal{I})$; $\mathcal{B} \leftarrow \operatorname{zeros}(h, w)$
 \diamond grassfire transformation
2: for $(i,j) \in \{0, \dots, h-1\} \times \{0, \dots, w-1\}$ do
3: if $\operatorname{color}(\mathcal{I}(i,j)) \notin \{\operatorname{black}, \operatorname{gray}\}$ then
4: $\mathcal{B}(i,j) \leftarrow 1 + \min\{\mathcal{B}(i-1,j), \mathcal{B}(i,j-1)\}$
5: for $(i,j) \in \{h-1, \dots, 0\} \times \{w-1, \dots, 0\}$ do
6: if $\operatorname{color}(\mathcal{I}(i,j)) \notin \{\operatorname{black}, \operatorname{gray}\}$ then
7: $\mathcal{B}(i,j) \leftarrow 1 + \min\{\mathcal{B}(i+1,j), \mathcal{B}(i,j+1)\}$
8: skeleton \leftarrow extract pixels making up the most intense
lines in the brightness map \mathcal{B}
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Generating Inspection Waypoints (2)

Input: \mathcal{I} : bitmap image; α : desired inspection quality, $0 < \alpha \leq 1$ Output: a set of inspection points

 \Diamond select inspection points

- 1: skeleton \leftarrow FILTER(skeleton)
- **2**: inspectionPts \leftarrow skeleton
- 3: currScore \leftarrow **VISSCORE**(\mathcal{I} , inspectionPts)
- 4: for $p \in \text{skeleton } \mathbf{do}$
- 5: newScore $\leftarrow VISSCORE(\mathcal{I}, inspectionPts \setminus \{p\})$
- 6: **if** newScore $\geq \alpha \lor \text{currScore} = \text{newScore}$ **then**
- 7: inspectionPts \leftarrow inspectionPts $\setminus \{p\}$
- 8: $\operatorname{currScore} \leftarrow \operatorname{newScore}$
- 9: return inspectionPts

Multi-Goal Motion Planning with Dynamics





Dynamics

- Express relation between input controls and resulting motions
- Necessary to plan motions that can be executed
- Impose significant challenges
- Constrain the feasible motions
- Often are nonlinear and highdimensional
- Give rise to nonholonomic systems
- State and control spaces are continuous
- Solution trajectories are often long



Computational complexity of motion planning with dynamics

- Point with Newtonian dynamics NP-Hard [DXCR 1993]
- Polygon Dubin's car Decidable [CPK 2008]
- General nonlinear dynamics Undecidable
 [Branicky 1995]

Dynamics

 Express relation between input controls and resulting motions



Modeled via physics-based engines

ROS/Gazebo, ODE, Bullet, PhysX general rigid-body dynamics friction and gravity

$$\dot{s} = f(s, u)$$

$$s = (x, y, \theta_0, v, \psi, \theta_1, \dots, \theta_n) \qquad u = (a, \omega)$$

$$\dot{x} = v \cos(\theta_0) \quad \dot{y} = v \sin(\theta_0) \quad \dot{\theta_0} = v \tan(\psi) \quad \dot{v} = a \quad \dot{\psi} = \omega$$

$$\dot{\theta_i} = \frac{v}{d} \left(\prod_{j=1}^{i-1} \cos(\theta_{j-1} - \theta_j) \right) (\sin(\theta_{i-1}) - \sin(\theta))$$

Introduce Discrete Layer to Guide the Search

Workspace decomposition provides

- discrete layer as adjacency graph G = (R,E)
- R denotes the regions of the decomposition
- E = {(ri,rj) | ri, rj in R are physically adjacent}

hcost(r) estimates the difficulty of reaching goal region from r

- defined as length of shortest path in G from r to goal [hcost(r1), hcost(r2),..., hcost(rn)]
- computed by running BFS/A* on G backwards from goal



cost heuristics over discrete layer guide search in A* fashion randomized sampling and PID controllers to expand motion tree

plans collision-free, dynamically-feasible, and low-cost solution trajectory



- Expand a tree T of collision-free and
- dynamically-feasible motions
- Select a state s from which to expand the tree
- Sample control input u
- ➤generate new trajectory by
- ≻applying u to s



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Complex Robot Models -Often Sets of Differential Equations



Figure 2: Vehicle models of a car, snake, and blimp used in the experiments.



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Guided Expansion of Motion Tree

Sampling-based motion planning

- generality: dynamics as black-box function s' = MOTION(s,u,dt)
- > continuous state/control spaces: probabilistic sampling to make it feasible
- high-dimensionality: search to find solution
- coupled with discrete abstractions
- provide simplified planning layer
- guide search in the continuous state/control spaces
- and motion controllers
- open up the black-box MOTION function
- facilitate search expansion

Formal guarantees

probabilistic completeness

Driven by



Architecture





Abstraction



Used to induce partition of motion tree into equivalence classes

 $V_i \equiv V_j \iff \text{TRAJ}(\mathcal{T}, V_i) \text{ provides same abstract information as } \text{TRAJ}(\mathcal{T}, V_j)$ region $(V_i) = \text{region}(V_j)$

 \Rightarrow equivalence class corresponding to abstract state $\langle r \rangle$

$$\Gamma_{\langle r \rangle} = \{ V : V \in \mathcal{T} \land \operatorname{region}(V) = r \}$$

 $\Rightarrow~$ partition of motion tree $\mathcal T$ into equivalence classes

$$\Gamma = \{ \Gamma_{\langle r \rangle} : |\Gamma_{\langle r \rangle}| > 0 \}$$

Graph Search for the Colored TSP

Let G = (V, E, color, cost) denote an undirected, colored, and weighted graph. Let $p_{\text{start}} \in V$ denote the start vertex. A sequence of vertices $\langle p_1, \ldots, p_r \rangle$ constitutes a **valid colored tour** if

- $\blacksquare \{p_1,\ldots,p_r\}=V,$
- $\blacksquare p_1 = p_{\text{start}},$
- $\forall i \in \{1, ..., r-1\} : (p_i, p_{i+1}) \in E$, and
- $\forall i, j, k \in \{1, ..., r\}$: black $\notin \{\operatorname{color}(p_i), \operatorname{color}(p_j), \operatorname{color}(p_k)\}$ and $i < j < k \land \operatorname{color}(p_i) \neq \operatorname{color}(p_j) \Longrightarrow \operatorname{color}(p_i) \neq \operatorname{color}(p_k)$.

An **optimal colored tour** is a colored tour with minimum cost, where the cost of the tour is defined as the sum of the weights associated with the edges of the tour.

Physical CTSP: integrate system dynamics such as angular change into cost function.

Rosin 2011 IJCAI – Best Paper, Morpion Solitaire with new Re MCS tree based on Complete Rollout and Recursive Search

- Not really MCTS, No Search Tree.
- In Each Level a Policy is Maintained, Updated and Refreshed state
- Updating Policy based on better Solutions Coming in from below
- Policy in Turn Influences the Rollouts
- Parameters: Level of Recursion, and Iteration Width
- Effective for TSPTW and many other Approaches
- Refinements: Beam / Diversity / Generalization





Learning Curve



Input: Iteration width (exploitation), nestedness (exploration) **Policy:** (city-to-city) mapping NxN -> IR to be learnt



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Playout

Algorithm 1 The playout algorithm		
1: playout (<i>state</i> , <i>policy</i>)		
2: sequence \leftarrow []		
3: while true do		
4: if <i>state</i> is terminal then		
5: return (score (<i>state</i>), <i>sequence</i>)		
6: end if		
7: $z \leftarrow 0.0$		
8: for m in possible moves for state do		
9: $z \leftarrow z + \exp(policy [code(m)])$		
10: end for		
11: choose a move with probability $\frac{exp(policy[code(move)])}{z}$		
12: $state \leftarrow play (state, move)$		
13: $sequence \leftarrow sequence + move$		
14: end while		

Adapt

Algorithm 2 The Adapt algorithm		
1:	Adapt (policy, sequence)	
2:	$polp \leftarrow policy$	
3:	$state \leftarrow root$	
4:	for move in sequence do	
5:	$polp [code(move)] \leftarrow polp [code(move)] + \alpha$	
6:	$z \leftarrow 0.0$	
7:	for m in possible moves for state do	
8:	$z \leftarrow z + \exp(policy [code(m)])$	
9:	end for	
10:	for m in possible moves for $state$ do	
11:	$polp [code(m)] \leftarrow polp [code(m)] - \alpha * \frac{exp(policy[code(m)])}{z}$	
12:	end for	
13:	$state \leftarrow play (state, move)$	
14:	end for	
15:	$policy \leftarrow polp$	

Search

Algorithm 3 The NRPA algorithm.		
1:	NRPA (level, policy)	
2:	if level == 0 then	
3:	return playout (root, <i>policy</i>)	
4:	else	
5:	$bestScore \leftarrow -\infty$	
6:	for N iterations do	
7:	$(result, new) \leftarrow NRPA(level - 1, policy)$	
8:	if result \geq bestScore then	
9:	$bestScore \leftarrow result$	
10:	$seq \leftarrow new$	
11:	end if	
12:	policy \leftarrow Adapt (policy, seq)	
13:	end for	
14:	return (bestScore, seq)	
15:	end if	
Theory...

The probability p_{ik} of choosing the move m_{ik} in a playout is the softmax function:

$$p_{ik} = \frac{e^{w_{ik}}}{\sum_j e^{w_{ij}}}$$

The cross-entropy loss for learning to play move m_{ib} is $C_i = -log(p_{ib})$. In order to apply the gradient we calculate the partial derivative of the loss: $\frac{\delta C_i}{\delta p_{ib}} = -\frac{1}{p_{ib}}$. We then calculate the partial derivative of the softmax with respect to the weights:

$$\frac{\delta p_{ib}}{\delta w_{ij}} = p_{ib}(\delta_{bj} - p_{ij})$$

Where $\delta_{bj} = 1$ if b = j and 0 otherwise. Thus the gradient is:

$$\nabla w_{ij} = \frac{\delta C_i}{\delta p_{ib}} \frac{\delta p_{ib}}{\delta w_{ij}} = -\frac{1}{p_{ib}} p_{ib} (\delta_{bj} - p_{ij}) = p_{ij} - \delta_{bj}$$

If we use α as a learning rate we update the weights with:

$$w_{ij} = w_{ij} - \alpha(p_{ij} - \delta_{bj})$$

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Praxis...

https://nms.kcl.ac.uk/stefan.edelkamp/lectures/pi1/programs/VRP.java

🗇 BlueJ: Konsole - Einführung CV	-		×	: Mill.Pair	
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Some Results Multi-Goal



Branch and Bound Search vs. Precurser LTLSyslop

More Results Multi-Goal

Branch and Bound Search with various heuristics and Monte-Carlo Tree Search



Inspection Benchmarks





map 19

map 40



map 45



map 35



map 61

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Results Inspection



Intermediate Summary Multi-Goal Task-Motion Planning

- Approach makes it possible to consider
 - ⇒ high-dimensional robotic systems with nonlinear dynamics and nonholonomic constraints
 - \Rightarrow visit all goal regions fast in suitable cost-minimizing order
 - \Rightarrow unstructured, complex environments

and efficiently computes

- $\Rightarrow\,$ collision-free, dynamically-feasible, low-cost trajectories that enable the robot to satisfy the task specification ϕ
- Offers probabilistic completeness

Temporal Planning: Time Does Matter

In general, activities have varying durations:

- Loading a package onto a truck is much quicker than driving the truck;
- Drinking a cup of tea takes longer than making it;
- Procrastinating tasks takes longer than doing them

Example: Zeno Domain

Initial State

Goal State



```
(define (domain zeno-travel)
(:requirements :durative-actions :typing :fluents)
(:types aircraft person city)
(:predicates (at ?x - (either person aircraft) ?c - city)
             (in ?p - person ?a - aircraft))
(:functions (fuel ?a - aircraft) (distance ?c1 - city ?c2 - city)
            (slow-speed ?a - aircraft) (fast-speed ?a - aircraft)
            (slow-burn ?a - aircraft) (fast-burn ?a - aircraft)
            (capacity ?a - aircraft) (refuel-rate ?a - aircraft)
            (total-fuel-used) (boarding-time) (debarking-time))
(:durative-action board
 :parameters (?p - person ?a - aircraft ?c - city)
 :duration (= ?duration boarding-time)
 :condition (and (at start (at ?p ?c))
                 (over all (at ?a ?c)))
 :effect (and (at start (not (at ?p ?c)))
              (at end (in ?p ?a))))
[...]
(:durative-action zoom
 :parameters (?a - aircraft ?c1 ?c2 - city)
 :duration (= ?duration (/ (distance ?c1 ?c2) (fast-speed ?a)))
 :condition (and (at start (at ?a ?c1))
                 (at start (>= (fuel ?a) (* (distance ?c1 ?c2) (fast-burn ?a)))))
 :effect (and (at start (not (at ?a ?c1)))
              (at end (at ?a ?c2))
              (at end (increase total-fuel-used
                         (* (distance ?c1 ?c2) (fast-burn ?a))))
              (at end (decrease (fuel ?a)
                         (* (distance ?c1 ?c2) (fast-burn ?a))))))
```

Zeno PDDL domain File

```
(define (problem zeno-travel-1)
                          (:domain zeno-travel)
                          (:objects plane - aircraft
                                   ernie scott dan - person
                                   city-a city-b city-c city-d - city)
                          (:init (= total-fuel-used 0) (= debarking-time 20) (= boarding-time 30)
Zeno
                                (= (distance city-a city-b) 600) (= (distance city-b city-a) 600)
                                (= (distance city-b city-c) 800) (= (distance city-c city-b) 800)
PDDL
                                (= (distance city-a city-c) 1000) (= (distance city-c city-a) 1000)
                                (= (distance city-c city-d) 1000) (= (distance city-d city-c) 1000)
Problem
                                (= (fast-speed plane) (/ 600 60)) (= (slow-speed plane) (/ 400 60))
                                (= (fuel plane) 750)
                                                          (= (capacity plane) 750)
                                (= (fast-burn plane) (/ 1 2)) (= (slow-burn plane) (/ 1 3))
FILE
                                (= (refuel-rate plane) (/ 750 60))
                                (at plane city-a) (at scott city-a) (at dan city-c) (at ernie city-c))
                          (:goal (and (at dan city-a) (at ernie city-d) (at scott city-d)))
                          (:metric minimize total-time)
```

Sequential and TemPORAL PLAN

- 0: (zoom plane city-a city-c) [100]
- 100: (board dan plane city-c) [30]
- 130: (board ernie plane city-c) [30]
- 160: (refuel plane city-c) [40]
- 200: (zoom plane city-c city-a) [100]
- 300: (debark dan plane city-a) [20]
- 320: (board scott plane city-a) [30]
- 350: (refuel plane city-a) [40]
- 390: (zoom plane city-a city-c) [100]
- 490: (refuel plane city-c) [40]
- 530: (zoom plane city-c city-d) [100]
- 630: (debark ernie plane city-d) [20]
- 650: (debark scott plane city-d) [20]

- 0: (zoom plane city-a city-c) [100]
- 100: (board dan plane city-c) [30] (board ernie plane city-c) [30]
- 100: (refuel plane city-c) [40]
- 140: (zoom plane city-c city-a) [100]
- 240: (debark dan plane city-a) [20] (board scott plane city-a) [30] (refuel plane city-a) [40]
- 280: (zoom plane city-a city-c) [100]
- 380: (refuel plane city-c) [40]
- 420: (zoom plane city-c city-d) [100]
- 520: (debark ernie plane city-d) [20] (debark scott plane city-d) [20]

SNAG: Sequential Plan TIME VS. Parallel Plan TIME

```
(zoom city-a city-c plane), (board dan plane city-c),
(refuel plane city-c), (zoom city-c city-a plane),
(board scott plane city-a), (debark dan plane city-a), (refuel plane city-a),
```

and

(board scott plane city-a), (zoom city-a city-c plane), (board dan plane city-c), (refuel plane city-c), (zoom city-c city-a plane), (debark dan plane city-a), (refuel plane city-a)

Different Plan OBJECTIVES

Fuel

Time

- 0: (board scott plane city-a) [30]
- 30: (fly plane city-a city-c) [150]
- 180: (board ernie plane city-c) [30] (board dan plane city-c) [30]
- 210: (fly plane city-c city-a) [150]
- 360: (debark dan plane city-a) [20] (refuel plane city-a) [53.33]
- 413.33: (fly plane city-a city-c) [150]
- 563.33: (fly plane city-c city-d) [150]
- 713.33: (debark ernie plane city-d)[20]
 - (debark scott plane city-d)[20]

- 0: (zoom plane city-a city-c) [100]
- 100: (board dan plane city-c) [30]
 - (board ernie plane city-c) [30]
 - (refuel plane city-c) [40]
- 140: (zoom plane city-c city-a) [100]
- 240: (debark dan plane city-a) [20]
 - (board scott plane city-a) [30]
 - (refuel plane city-a) [40]
- 280: (fly plane city-a city-c) [150]
- 430: (fly plane city-c city-d) [150]
- 580: (debark ernie plane city-d) [20]
 - (debark scott plane city-d) [20]

Durative Actions?



Durative Actions?



Durative Actions in PDDL 2.1



at start at end

PDDL Example (i)

- . re-action LOAD-TRUCK
- :parameters
- (?obj obj ?truck truck ?loc -
- :precondition (= ?duration 2)
- :cond
- (and .11 (at ?truck
- (at start (at ?obj ?loc)))
- :eff----
- (ar art (not (at ?obj
- (at end (in ?obj ?truck)))

Beware of self-overlapping actions!

PDDL Example (ii)

- (:durative-action open-barrier
- :parameters
- (?loc location ?p person)
- :duration (= ?duration 1)
- :condition
- (and (at start (at ?loc ?p)))
- :effect
- (and (at start (door-open ?loc))
- (at end (not (door-open ?loc))))

PDDL Example (ii)

- (:durative-actio
- :parameters
- (?loc locatio
- :duration (= ?
- :condition
- (and (at st
- :effect



- (and (at start (door-open ?loc))
- (at end (not (door-open ?loc))))

Durative Actions

(Fox and Long, ICAPS 2003)



Planning with Snap Actions (i)



Challenge 1: What if B interferes with the goal?
@PDDL 2.1 semantics: no actions can be executing in a goal state.

Solution: add ¬As, ¬Bs, ¬Cs.... to the goal (or make this implicit in a temporal planner.)

Planning with Snap Actions (ii)



Challenge 2: what about **over all** conditions?

If A is executing, inv_A must hold.

Solution:

In every state where As is true: inv_A must also be true

Or: (imply (As) inv_A)

Violating an invariant then leads to a **dead-end**.

Planning with Snap Actions (iii)



- Challenge 3: where did the durations go?
 More generally, what are the temporal constraints?
- -Logically sound ≠ temporally sound.

Option 1: Decision Epoch Planning Term from Cushing et al, IJCAI 2007

- Search with time-stamped states and a priority queue of pending end snap-actions.
- See Temporal Fast Downward (Eyerich, Mattmüller and Röger, ICAPS 2009); Sapa (Do and Kambhampati, JAIR 2003), and others.
- > In a state S, at time t and with queue Q, either:
- > Apply a start snap-action \dot{A} (at time t)
- > Insert A _ into Q at time (t + dur(A))
- > S'.t = S.t + ε
- Remove and apply the first end snap-action from Q.
- > S'.t set to the scheduled time of this, plus ε

Running through our example...



Decision Epoch Planning: The snag

- Must fix start- and end-timestamps at the point when the action is started.
- -Used for the priority queue



OPTION 2: Simple Temporal Networks

"Planning with Problems Requiring Temporal Coordination." A. I. Coles, M. Fox, D. Long, and A. J. Smith. AAAI 08. https://local.cis.strath.ac.uk/research/publications/papers/strath_cis_publication_2248.pdf "Managing concurrency in temporal planning using planner-scheduler interaction." A. I. Coles, M. Fox, K. Halsey, D. Long, and A. J. Smith. Artificial Intelligence. 173 (1). 2009.

a Simple Temporal Problem?

All our constraints are of the form:

- \succ ε ≤ t(i+1) t(i) (c.f. sequence constraints)
- → $dur_{min}(A) \le t(A) t(A) \le dur_{max}(A)$

> Or, more generally, $lb \le t(j) - t(i) \le ub$

- Is a Simple Temporal Problem
- "Temporal Constraint Networks", Dechter, Meiri and Pearl, AIJ, 1991
- Good news is polynomial
 Bad news in planning, we need to solve it a lot....

Example

- John travels to work either by car (30-40 min) or by bus (>= 60 min)
- Fred travels to work either by car (20-30 min) or in a carpool (40-50 min)
- Today John left between 7:10 and 7:30am.
- Fred arrived at work between 8:00 and 8:10am.
- John arrived at work 10-20min after Fred left home.

Visualize TCSP as Directed Constraint Graph



Simple Temporal Network



Simple Temporal Network:

A set of time points X_i at which events occur. Unary constraints $(a_0 \le X_i \le b_0)$ or $(a_1 \le X_i \le b_1)$ or ... Binary constraints $(a_0 \le X_i - X_i \le b_0)$ or $(a_1 \le X_i - X_i \le b_1)$ or ...

STN



Shostak (1981) A simple temporal problem is consistent if and only if the distance graph has no cycles.

→ The consistency and the minimal network of an STP can be determined in cubic time using all-pairs shortest path search.

To Query STN Map to Distance Graph G_d

Edge encodes an upper bound on distance to target from source.



Shortest Paths of G_d



d-graph
STN Minimum Network

	0	1	2	3	4		0	1	2	3	4
0	0	20	50	30	70	0	[0]	[10,20]	[40,50]	[20,30]	[60,70]
1	-10	0	40	20	60	1	[-20,-10]	[0]	[30,40]	[10,20]	[50,60]
2	-40	-30	0	-10	30	2	[-50,-40]	[-40,-30]	[0]	[-20,-10]	[20,30]
3	-20	-10	20	0	50	3	[-30,-20]	[-20,-10]	[10,20]	[0]	[40,50]
4	-60	-50	-20	-40	0	4	[-70,-60]	[-60,-50]	[-30,-20]	[-50,-40]	[0]

d-graph

STN minimum network

Test Consistency: No Negative Cycles



d-graph

Latest Solution

Node 0 is the reference.



4

d-graph

Earliest Solution

Node 0 is the reference.

						<u> </u>
	0	1	2	3	4	
0	0	20	50	30	70	-10 -30 -10
1	-10	0	40	20	60	
2	-40	-30	0	-10	30	-40
3	-20	-10	20	0	50	-60
4	-60	-50	-20	-40	0	
						70

d-graph

Feasible Values

	0	1	2	3	4
0	0	20	50	30	70
1	-10	0	40	20	60
2	-40	-30	0	-10	30
3	-20	-10	20	0	50
4	-60	-50	-20	-40	0

- X₁ in [10, 20]
- X₂ in [40, 50]
- X₃ in [20, 30]
- X₄ in [60, 70]

d-graph

Back to Planning: Latest possible times? (Maximum Separation)

$$t(A) - t(Z) \le 4$$

 $t(B) - t(Z) \le 8$
 $t(C) - t(Z) \le 10$

('A comes no more than 4 time units after Z')



Latest possible times?

$$t(A) - t(Z) \le 4$$

$$t(B) - t(Z) \le 8$$

$$t(C) - t(Z) \le 10$$
('B comes no
more than 2 time

$$t(C) - t(B) \le 1$$
('B comes no
more than 2 time
units after A')



Earliest possible times? (Minimum Separation)

- For latest possible time: find the **shortest path**
- For earliest possible times...?

Earliest possible times?

$$2 \le t(A) - t(Z)$$

 $4 \le t(B) - t(Z)$

$$3 \le t(B) - t(A)$$

 $1 \le t(C) - t(B)$



Hacking algorithms

- > Longest path from Z to C?
- > = Shortest **negative** path from C to Z

 $2 \le t(A) - t(Z)$

Multiply both sides by -1: -2 > -t(A) + t(Z)

 $p \ge q$ is the same as $q \le p$: - t(A) + t(Z) < -2

> Rearrange LHS: t(Z) - t(A) < -2

Earliest possible times?

- $2 \le t(A) t(Z) -2 \ge -t(A) + t(Z) -t(A) \le -2$ $4 \le t(B) - t(Z) -4 \ge -t(B) + t(Z) -t(B) \le -4$
- $\begin{array}{lll} 3 <= t(B) t(A) & -3 >= -t(B) + t(A) & t(A) t(B) <= -3 \\ 1 <= t(C) t(B) & -1 >= -t(C) + t(B) & t(B) t(C) <= -1 \end{array}$



Simple Temporal Networks (i)

- Can map STPs to an equivalent digraph:
- One vertex per time-point (and one for 'time zero');
- ► For $lb \le t(j) t(i) \le ub$:
- > An edge (i \rightarrow j) with weight *ub*.
- > An edge $(j \rightarrow i)$, with weight *-lb*
- > (c.f. $lb \le t(j) t(i) \longrightarrow t(j) t(i) \le -lb$)

Example STN



Simple Temporal Networks (ii)

 Solve the shortest path problem (e.g. using Bellman-Ford) from/to zero

 $-dist(0,j)=x \rightarrow$ maximum timestamp of j = x

 $-dist(j,0)=y \rightarrow minimum timestamp of j = -y$

• If we find a **negative cycle** then the temporal constraints are inconsistent:



"Incremental Constraint-Posting Algorithms in Interleaved Planning and Scheduling." A. J. Coles, A. I. Coles, M. Fox, and D. Long. Proceedings of the Workshop on Constraint Satisfaction Techniques for Planning and Scheduling, ICAPS09 <u>https://local.cis.strath.ac.uk/research/publications/papers/strath_cis_publication_2409.pdf</u>



STN Simplifies For Partially Ordered Plans



Transitively implied edges omitted for clarity:
 –e.g. all the drive/board ends before work end;

-All the drive/board starts after work start.

Public Transport Example

- Drivers have working hours;
- Bus routes have fixed durations and start and end locations.
- Goals are that each bus route is done.
- The routes have timetables that they must follow.

Temporal Planning: Public Transport



Actions have:

- Conditions and Effects at the start and at the end;
- Invariant/overall conditions;
- Durations constraints:

```
(= ?duration 4)
(and (>= ?duration 2) (<= ?duration 4))</pre>
```

"Planning with Problems Requiring Temporal Coordination." A. I. Coles, M. Fox, D. Long, and A. J. Smith. AAAI 2008. "Managing concurrency in temporal planning using planner-scheduler interaction." A. I. Coles, M. Fox, K. Halsey, D. Long, and A. J. Smith. Artificial Intelligence. 173 (1) (2009).

Planning with Snap Actions



Three Challenges:

- Make sure ends can't be applied unless starts have.
- Overall Conditions.
- Duration constraints.

"Planning with Problems Requiring Temporal Coordination." A. I. Coles, M. Fox, D. Long, and A. J. Smith. AAAI 2008. "Managing concurrency in temporal planning using planner-scheduler interaction." A. I. Coles, M. Fox, K. Halsey, D. Long, and A. J. Smith. Artificial Intelligence. 173 (1) (2009).

Planning with Snap Actions and STNs



Constraints:

- $W_{+} W_{+} \ge 2$ $W_{+} W_{+} \le 4$ $R1_{+} \ge W_{+} + \varepsilon$ $R1_{+} R1_{+} = 2$ $R3_{+} \ge R1_{+} + \varepsilon$
- $R3_{-}$ $R3_{-}$ = 3
- $W_{\rightarrow} >= R3_{\rightarrow} + \epsilon$

"Planning with Problems Requiring Temporal Coordination." A. I. Coles, M. Fox, D. Long, and A. J. Smith. AAAI 2008. "Managing concurrency in temporal planning using planner-scheduler interaction." A. I. Coles, M. Fox, K. Halsey, D. Long, and A. J. Smith. Artificial Intelligence. 173 (1) 2009.

Temporal Task-Motion Planning

- Time is money
- Real-world has and needs time constraints



- Combining task with motion planning "holy grail" in robotics
- Multiple goals is planning for longer-term plans in form of tours
- Inspection problems can be solved via waypoint finding
- Robots have complex, non-linear dynamics



Temporal Task-Motion Planning



Crucial for Robotics, Logistics, Surgery, VR, ...

No (convincing) solution so far

High-Level Task, Low-Level Motion Planning

Solutions in Discrete World only Approximate

→ Replanning needed.



Randomized Roadmaps



PDDL Task Planning and TILS



TIL = Timed Initial Literal

(at timepoint (fact)) (at timepoint (not fact))

Specified in initial state

Leads to time windows for actions

Interface with PDDL Temporal Planner (OPTIC)

Input

(at auv v0) (connected v0 v1) (connected v0 v2) (connected v1 v2) (= (traveltime v0 v1) 0.8) (= (traveltime v0 v2) 1.5) (= (traveltime v1 v2) 0.7) (located task1 v1) (located task2 v2) (at 1.1 (tw_open task1)) (at 2.1 (not (tw_open task1))) (at 2.3 (tw_open task2)) (at 3.3 (not (tw_open task2)))

Output



Interface with Specialised Solvers (BnB, MCS, Random)

- SSSP-Reduced Roadmap Graph via Dijkstra Calls (Computed Prior to the Search)
- Cities = Waypoints, Current Position of Robot = Depot
- Open Tour
- Pairwise SSSP Distance Table
- Time Windows



B&B tree (/home/temp/test.bak Os)

Example











Pickup and Deliveries



Framework

Input: multi-goal motion planning problem with time windows, pickups, deliveries, capacities

Output: collision-free and dynamically-feasible trajectory ζ that seeks to maximize ${}_{\rm GOALS}(\zeta)$

- 1: $RM \leftarrow \text{ConstructRoadmap}(\mathcal{O}, \mathcal{G})$
- 2: $\Xi \leftarrow \text{SHORTESTPATHS}(RM, \mathcal{G})$
- 3: $\mathcal{T} \leftarrow \text{INITIALIZEMOTIONTREE}(s_{\text{init}})$
- 4: $\mathcal{X} \leftarrow \text{INITIALIZEEQUIVALENCECLASSES}(s_{\text{init}})$
- 5: while $TIME() < t_{max}$ and not solved and not converged do
- 6: $\mathcal{X}_{\text{key}} \leftarrow \text{SELECTEQUIVALENCECLASS}(\mathcal{X})$
- 7: $\mathcal{X}_{\text{key}}.\sigma \leftarrow \text{DISCRETESOLVER}(RM, \Xi, \text{key})$
- 8: EXPANDMOTIONTREE($\mathcal{T}, \mathcal{X}_{key}.\sigma$)
- 9: UPDATEEQUIVALENCECLASSES(\mathcal{X})

10: **return** trajectory ζ in \mathcal{T} that maximizes $GOALS(\zeta)$

Integration with Specialized Solvers



Interface with PDDL

Input

```
(:durative-action execute_task_pickup
:parameters (?v - vehicle ?wp - waypoint ?t - task)
:duration ( = ?duration (taskduration ?t))
 :condition (and
   (at start (at ?v ?wp)) (at start (located ?t ?wp)
   (at start (todo ?t))
   (at start (<= (+ (customer ?wp) (cap ?v)) (max_cap ?v)))</pre>
   (at start (is-pickup ?wp)) (at start(tw_open ?t)))
:effect (and
   (at start (not (todo ?t))) (at end (visited ?wp))
   (at end (increase (cap ?v) (customer ?wp)))
   (at end (decrease (profit ?v) (customer ?wp)))
   (at end (completed ?t)))
(:durative-action execute_task_delivery
:parameters (?v - vehicle ?wp1 ?wp2 - waypoint ?t - task)
 :duration ( = ?duration (taskduration ?t))ov
 :condition (and
   (at start (at ?v ?wp1)) (at start (located ?t ?wp1))
   (at start (todo ?t)) (at start (is-delivery ?wp1))
   (at start (and (visited ?wp2) (link ?wp2 ?wp1)))
   (at start (tw_open ?t)))
:effect (and
   (at start (not (todo ?t))) (at end (visited ?wp1))
   (at end (increase (cap ?v) (customer ?wp1)))
   (at end (decrease (profit ?v) (customer ?wp1)))
   (at end (completed ?t))))
```

Output

v1 v3 v5 v4 v2 0.0 1.26 3.22 12.55 21.11

coffee_errors.pddl					
1	(define COFFEE				
2					
3	(requirements				
4	:typing)				
5					
6	(:types room - location				
7	robot human _ agent				
8	furniture door - (at ?l - location)				
9	kettle ?coffee cup water - movable				
10	location agent movable - object)				
11					
12	(:predicates (at ?l - location ??o - object)				
13	(have ?m - movable ?a - agent)				
14	<pre>(hot ?m - movable) = true</pre>				
15	<pre>(on ?f - furniture ?m - movable))</pre>				
16					
17	(:action boil				
18	<pre>:parameters (?m - movable \$k - kettle ?a - agent)</pre>				
19	:preconditions (have ?m ?a)				
20	:effect (hot ?m))				
21					
Line 20,	Column 22 Spaces: 2 PDDL				

Scenes (With Randomized Road Maps)

Snake Model

Car Model





(a) MC (B) OPtic (C) Random



Adjusting Time Windows $[(1-\epsilon)t_i, (1+\epsilon)t_i]$



Adjusting Vehicle Capacity



RunTime Distribution

Bottom to Top: RoadMap Constr., APSP, Collision & Simulate, Discr. Solver, Other


Conclusion

Full-Fledged Solution:

- high-dimensional robotic systems with nonlinear dynamics and
- nonholonomic constraints
- > visit all goal regions fast in suitable cost-minimizing order
- > unstructured, complex environments

and efficiently computes

- > collision-free, dynamically-feasible, low-cost trajectories that
- enable the robot to satisfy the CPSPTW+PD task specification

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