# Adversarial Search 

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Games, man vs. algorithm

- Deep Blue
- Alpha Go
- Deep Stack
- Why Games, actually?

Notes
Please note, the hyperlinks at the main slides are not active in the slides with notes. Hyperlinks within the notes should be active, though.

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- Why Games, actually?

Games are interesting for AI because they are hard (to solve).

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## More: Adversarial Learning



## Video: Adversing visual segmentation

Vision for Robotics and Autonomous Systems, http://cyber.felk.cvut.cz/vras

- Fooling Tesla autopilot by adversarial attack:

Elements of the game

- $s_{0}$ : The initial state


Notes
Defining a game as a kind of search problem:
Considering the notation, we are making slight transition from [1] to [2].

- Players: $P=\{1,2, \ldots, N\}$ (often just $N=2$ )
- Transition functions: $S \times A \rightarrow S$.
- Terminal utilities: $S \times P \rightarrow R .(R-$ as a Reward $)$

What are we loking for? A strategy/policy $S \rightarrow A$

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- terminal-test(s). Game over?
- terminal-Utility $(s, p)$. What is prize? Examples for some
 games ...
https://commons.wikimedia.org/wiki/File:
Tic-tac-toe_5.png


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- Zero-sum: players have opposite utilities (values)
- Zero-sum: playing against opponent

Notes
Most common games-such as chess-have these properties:

- two-player
- turn-taking
- deterministic with perfect information (a.k.a. deterministic, fully observable environments)

In some games, there is imperfect information (evironment is not fully observable). E.g., poker - no access to what cards opponents hold.

- Zero-sum: players have opposite utilities (values)
- Zero-sum: playing against opponent
- General game: independent utilities
- General game: cooperations, competition, ...

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## Game Tree(s)



Terminal-Utility $(s, \mathbf{x})$

Init state, ACTIONS function, and RESULT function defines game tree.

Note: game tree as opposed to search tree. Game tree are all possible evolutions of the game. (With standard search, we similarly had state space graph vs. search tree.)

State Value $V(s)$
$V(s)$ - value $V$ of a state $s$ : The best utility achievable from this state.

$$
V(s)=\max _{s^{\prime} \in \operatorname{children}(s)} V\left(s^{\prime}\right)
$$

Notes
Think about the State Value. It is a theoretical construct, definition. Depending on the problem, there may be various computational algorithms.
In a game, what State Values are known? Usually, only terminal states.
Think, for a moment, you are the only player. You can control every step. How would you compute the $V(s)$ for a given state $s$ ?

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Two-ply game: max for me, min for the opponent.

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## Notes

One move consists of two plies (half-moves).
I'm the player that starts (state A) and want to decide what to play; actions/plies $a_{1}, a_{2}, a_{3}$ are the options. B, C, D are the possible outcomes of my moves (plies). Now the opponent is about to play. The numbers in terminal states denote my profit/utility.
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MAX (x)


TERMINAL
Utility

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Max step: I want to maximize my outcome.
Min step: Opponent wants to maximize his outcome which is equivalent to minimizing my outcome.

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## Notes

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## Minimax algorithm

function MINIMAX(state) returns an action
function MIN-VALUE(state) returns a utility value $v$
function MAX-VALUE(state) returns a utility value $v$

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function MINIMAX(state) returns an action
return argmax MIN-VALUE(RESULT(state, a)) $a \in$ Actions(s)
end function
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function MIN-VALUE(state) returns a utility value $v$
if TERMINAL-TEST(state) then return UTILITY(state) end if
$v \leftarrow \infty$
for all ACTIONS(state) do
$v \leftarrow \boldsymbol{\operatorname { m i n }}(v, \operatorname{MAX}-\operatorname{VALUE}(\operatorname{RESULT}($ state,$a)))$
end for
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if TERMINAL-TEST(state) then return UTILITY(state)
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$v \leftarrow-\infty$
for all ACTIONS(state) do
$v \leftarrow \boldsymbol{\operatorname { m a x }}(v, \operatorname{MIN}-\operatorname{VALUE}(\operatorname{RESULT}($ state,$a)))$
end for
end function

A two ply game, down to terminal and back again ...
function minimax $(s)$ returns a $\operatorname{argmax} \operatorname{MINVAL}(\operatorname{RES}(s, a))$ MAX $a \in$ Actions(s)
end function
function MINVAL(s) returns $v$
if $\operatorname{terminal}(s)$ then $\operatorname{util}(s)$
end if
$v \leftarrow \infty$
for all ACTIONS(s) do
$v \leftarrow \boldsymbol{\operatorname { m i n }}(v, \operatorname{MAXVAL}(\operatorname{RES}(s, a)))$
end for

end function

```
function MAXVAL(s) returns \(v\)
    if TERMINAL( \(s\) ) then UTil( \(s\) )
    end if
    \(v \leftarrow-\infty\)
    for all ACTIONS(s) do
        \(v \leftarrow \max (v, \operatorname{MiNVAL}(\operatorname{RES}(s, a)))\)
        end for
    end function
```

Before going to the animation on the next slide, try to follow the algorithm by a pencil and paper.

A two ply game, recursive run

Efficiency/complexity:

- Exhaustive DFS
- Time $O\left(b^{m}\right)$
- Space $O(b m)$

Chess $b \approx 35, m \approx 100 \ldots$

- We cannot go(dive) to the end
- Can we save something?

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Can we do better? How?

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Nodes (sub-trees) worth visiting

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Notes
Constraining the possible node values as search progresses...

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## $\alpha-\beta$ pruning

$\alpha$ highest (best) value choice found so far for any choice along MAX $\beta$ lowest (best) value choice found so far for any choice along MIN

$v$ value of the state

Notes
Functions scope: MAX-VALUE MIN-VALUE. The terminal nodes are served/answered within the MAX-VALUE function.
Once a node (subtree) is exhausted (fully expanded), values propage to towards the root In MAX nodes $\alpha$ is changing and $\beta$ is stopping, in MIN nodes $\beta$ is changing and $\alpha$ is stopping.

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$\beta$ lowest (best) value choice found so far for any choice along MIN


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original: Time: $O\left(b^{m}\right)$

- how to select nodes?
- perfect ordering?


It is clear that ordering of child nodes matters. Draw tree of $\alpha-\beta$ search in case of perferct ordering. Effective branching factor becomes $\sqrt{b}$ instead of $b$ which effectively doubles the depth can be searched: Time: $O\left(b^{m / 2}\right)$

# function ALPHA-BETA-SEARCH(state) returns an action 

$v \leftarrow \operatorname{MAX}-\operatorname{VALUE}($ state, $\alpha=-\infty, \beta=\infty$ )
return action corresponding to $v$
end function

Notes
Take the tree from the previous slide and try to go step-by-step, watch $\alpha, \beta$ and $v$
function ALPHA-BETA-SEARCH(state) returns an action $v \leftarrow$ MAX-VALUE(state, $\alpha=-\infty, \beta=\infty$ ) return action corresponding to $v$
end function
function MAX-VALUE(state, $\alpha, \beta$ ) returns a utility value $v$
if TERMINAL-TEST(state) return UTILITY(state)
$v \leftarrow-\infty$
for all ACTIONS(state) do
$v \leftarrow \boldsymbol{\operatorname { m a x }}(v, \operatorname{MIN}-\operatorname{VALUE}(\operatorname{RESULT}($ state,$a), \alpha, \beta))$
if $v \geq \beta$ return $v$
$\alpha \leftarrow \max (\alpha, v)$ end for
end function

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end for
end function
function MIN-VALUE(state, $\alpha, \beta$ ) returns a utility value $v$
if TERMINAL-TEST(state) return UTILITY(state)
$v \leftarrow \infty$
for all ACTIONS(state) do
$v \leftarrow \boldsymbol{\operatorname { m i n }}(v, \operatorname{MAX}-\operatorname{VALUE}(\operatorname{RESULT}($ state,$a), \alpha, \beta))$
if $v \leq \alpha$ return $v$ $\beta \leftarrow \min (\beta, v)$

## end for

Take the tree from the previous slide and try to go step-by-step, watch $\alpha, \beta$ and $v$

Recall: Iterative deepening DFS (ID-DFS)

- Start with maxdepth = 1
- Perform DFS with limited depth. Report success or failure.
- If failure, forget everything, increase maxdepth and repeat DFS

The "wasting" of resources is not too bad. Recall:

- Most nodes are at the deepest levels.
- Asymptotic complexity unchanged.


Bonus for $\alpha-\beta$ pruning: previous "shallower" iterations can be reused for node ordering.

## Notes

$\alpha-\beta$ pruning is good. Still, in chess, for example, there is no way we can compute till the end.
Time is limited. We need to respond within a certain amount of time.
Possible solution: iterative deepening search. If I can't complete the computation for the current depth, I can
use the previous shallower one that finished.

Imperfect but real-time decisions: iterative deepening

$$
\operatorname{H-MinimAx}(s, d)=
$$

Even with perfect ordering, $\alpha-\beta$ pruning does not save us.
One problem left: can't compute till then end, need to cut off, need for Evaluation function.

Imperfect but real-time decisions: iterative deepening

$$
\begin{aligned}
\operatorname{H-minimax}(s, d) & = \\
\operatorname{EVAL}(s) & \text { if } \operatorname{CUTOFF-TEST}(s, d)
\end{aligned}
$$

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Imperfect but real-time decisions: iterative deepening

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\begin{array}{rlll}
\operatorname{H-minimax}(s, d)= & & \\
\operatorname{EVAL}(s) & \text { if } & \operatorname{CUTOFF}-\operatorname{TESt}(s, d) \\
\max _{a \in \operatorname{ACTIONS}(s)} \operatorname{H-MinimAx}(\operatorname{RESULT}(s, a), d+1) & \text { if } & \operatorname{PLAYER}(s)=\max
\end{array}
$$

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$$<br>$$
\min _{a \in \operatorname{ACTIONS}(s)} \operatorname{H-MINIMAX}(\operatorname{RESULT}(s, a, d+1)) \text { if } \operatorname{PLAYER}(s)=\operatorname{MIN}
$$

Even with perfect ordering, $\alpha-\beta$ pruning does not save us.
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Cutting off search and evaluation functions

```
Replace
if TERMINAL-TEST(s) then return TERMINAL-UTILITY(s) with:
if CUTOFF-TEST( \(s, d\) ) then return EVAL( \(s\) )
```

Historical note: cutting search off earlier and use of heuristic evaluation functions proposed by Claude Shannon in Programming a Computer for Playing Chess (1950).

EVAL(s) - Evaluation functions
(estimate of) State value for non-terminal states
We need an easy-to-compute function correlated with "chance of winning". For chess:

- Material value for pieces-1 for pawn, 3 for knight/bishop, 5 for rook, 10 for queen. (minus opponent's pieces)
- Finetuning: 2 bishops are worth 6.5 ; knights are worth more in closed positions..
- Other features worth evaluating: controlling the center of the board, good pawn structure (no double pawns), king safety...


## Notes

For many problems it is not so easy to find/construct proper function. We may try more functions and combine them conveniently.

$$
f_{1}(s)=\text { number of white pawns - number of black pawns }
$$

How to tune weights $w_{i}$ ?
or Deep Nets! Yeah!
How to get training data for supervised learning? More later.

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$$
\operatorname{EvaL}(s)=w_{1} f_{1}(s)+w_{2} f_{2}(s)+\cdots w_{n} f_{n}(s)
$$

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EVAL(s) - Problems
What if something important happens just beyond the current horizon of the search?


Additional improvements:

- "Killer moves" -capturing opponent's pieces, check etc.-should be considered first.
- Quiescence search - EVAL function should be applied only once things calm down. During capturing of pieces, depth should be locally increased.

Computer play vs. grandmaster play

- Computers are better since 1997 (Deep Blue defeating Garry Kasparov).
- The way they play is still very different: "dumb", relying on "brute force".
- Grandmasters do not excel in being able to compute very deep-many moves ahead.
- They play based on experience: super-effective pruning and evaluation functions.
- They consider only 2 to 3 moves in most positions (branching factor).


## References

## Chapter 5, "Adversarial search" in [1].

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