Problem solving by search

Tomáš Svoboda and Matěj Hoffmann

Vision for Robots and Autonomous Systems, Center for Machine Perception
Department of Cybernetics
Faculty of Electrical Engineering, Czech Technical University in Prague

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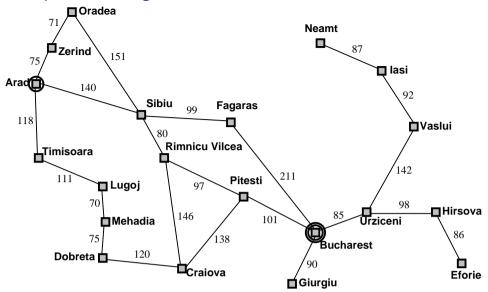
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Notes -

Outline

- ► Search problem.
- ► State space graphs.
- Search trees.
- ► Strategies: which tree branches to choose?
- Strategy/Algorithm properties.
- Programming infrastructure.

Example: Traveling in Romania



Notes -

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Ok, start with a simple one, almost everybody knows about the navigation - path planning problem. Waze, Garmin, . . .

Can you think about more problems?

For example:

- Touring problems. Special case: Traveling salesperson problem each city must be visited exactly once.
- VLSI (chip) layout.
- ..

Goal:

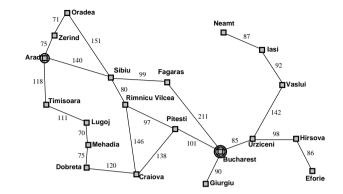
be in Bucharest

Problem formulation:

states: position in a city (cities actions: drive between cities

Solution

Sequence of cities (path) (action sequence [2])



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Notes -

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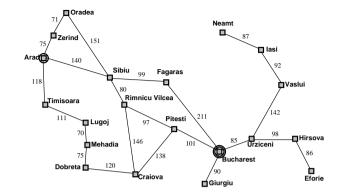
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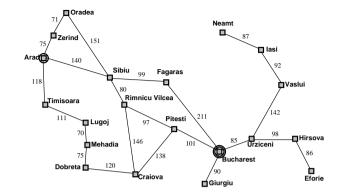
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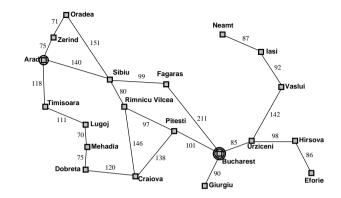
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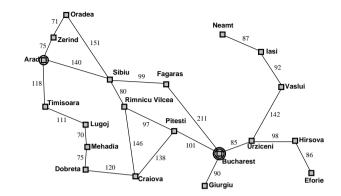
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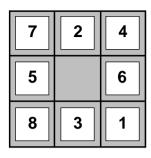
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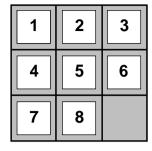


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Notes

Example: The 8-puzzle





Start State

Goal State

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states? actions? solution? cost?

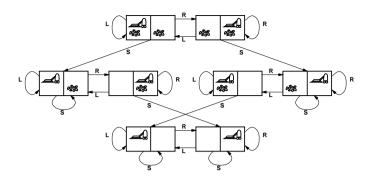
Notes -

Also known as n-1 puzzle.

- States: Location of each of the 8 tiles and the blank.
- Number of states: 9!
- Initial state: any state. (Note that any given goal state can be reached from exactly half of the initial states.)
- Actions: Movements of the blank space: Left, Right, Up, Down (or a subset of these)
- Solution / goal test: Check whether state matches the goal configuration.
- Path cost: nr. steps in the path (each step costs 1)

Toy problem (3.2.1) from [2].

Example: Vacuum cleaner



states? actions? solution? cost?

Notes

- States: Determined by agent location and dirt location. The agent is in one of two locations, each of which may or may not contain dirt.
- Number of states: 2×2^2 (two possible choices for agent location; for every location, choice dirt vs. no dirt). For *n* locations: $n \times 2^n$
- Initial state: any state
- Actions: Left, Right, Suck (larger envs. can have also Up and Down)
- Solution / goal test: Are all squares clean?
- Path cost: nr. steps in the path (each step costs 1)

Toy problem (3.2.1) from [2].

- ► State space (including Start/Initial state): position, board configuration,
- ► Actions: drive to, Up, Down, Left ...
- ▶ Transition model : Given state and action return state (and cost)
- ▶ Goal test : Are we done?

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We will use the terminology throught the next 5-6 lectures; also for Markov (Sequential) Decision Processes, Reinforcement Learning

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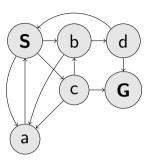
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State Space Graphs

State space graph: a representation of a search problem

- Graph Nodes states are abstracted world configurations
- Arcs represent action results
- ► Goal test a set of goal nodes

Each state occurs only once in a state (search) space.

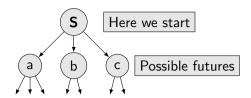


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Notes -

Formalizing a real world problem – (creating) a state space graph – could be a problem in itself. I put creating into brackets as it may be also infinite.

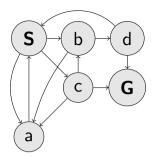
Search Trees



- ► A "what if" tree of plans and their outcomes
- Start node is the root
- ► Children are successors
- ▶ Nodes show/contains states, but correspond to *plans* that achieve those states

Notes -

- What if decision about an action, repeats . . .
- Nodes in the search tree are not the same as the nodes in the state space graph.



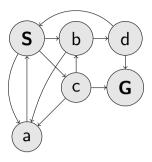
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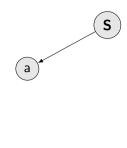
How big is the search tree?

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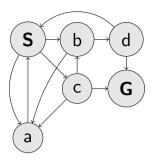


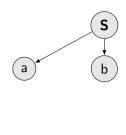


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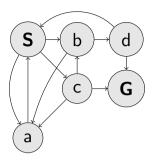


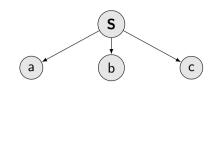


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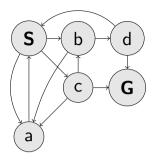


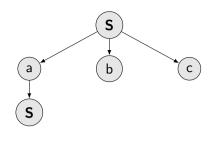


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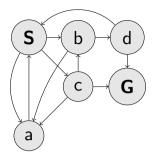


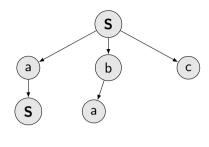


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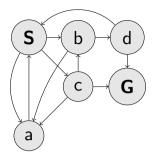


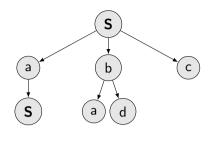


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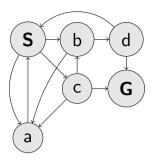


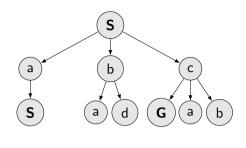


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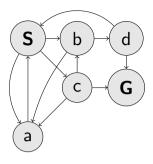


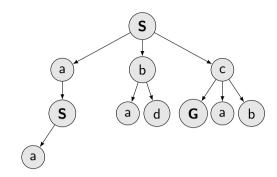


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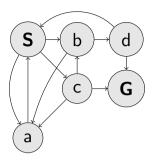


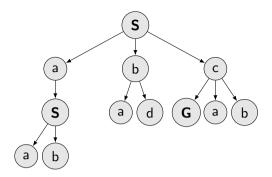


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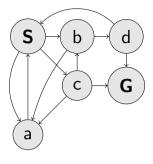


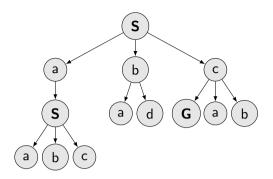


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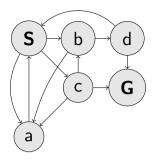


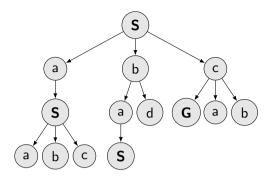


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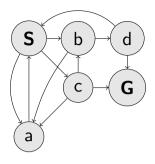


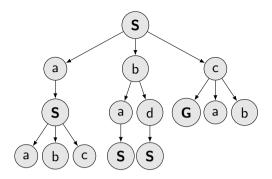


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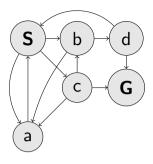


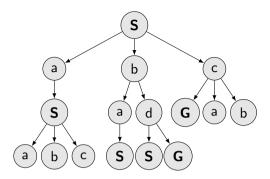


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From problem/transition graph to search tree (Romania) (a) The initial state □ Oradea Neamt (b) After expanding Arad Sibiu Sibiu Fagaras 118 Vaslui Timisoara Rimnicu Vilcea Pitesti (c) After expanding Sibiu Hirsova Mehadia

 ${\sf Problem/transition\ graph\ is\ revealed\ incrementally}.$

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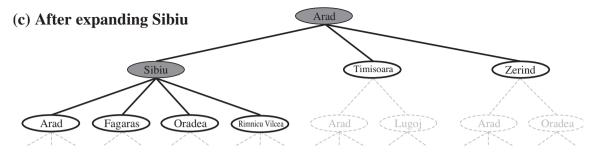
The revealing strategy can be visualized as a search tree.

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Eforie Arad Fagaras

Search elements - unvisited, dead, alive states



- Expand plans possible ways (tree nodes).
- ► Manage/Maintain fringe (or frontier) of plans under consideration.
- Expand new nodes wisely(?).



function TREE_SEARCH(problem) return a solution or failure

initialize by using the initial state of the problem loop

if no candidates for expansion then return failure

and if

if the node contains a goal state then return the solution

end if

Expand the node and add the resulting nodes to the tree

end loop

end function

Notes -

A *general* tree search algorithm. Individual search algorithms vary primarily in how they choose which state to expand next – the "search strategy".



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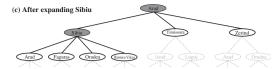
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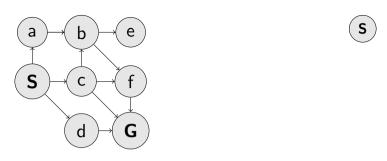
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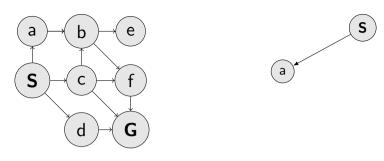
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Which nodes to *explore*? What are the properties of a strategy/algorithm?

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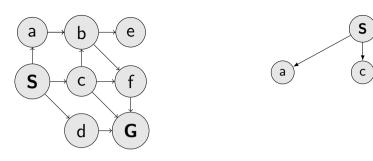
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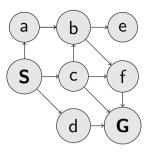
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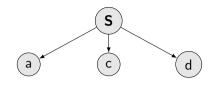


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Notes -

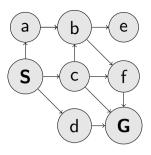


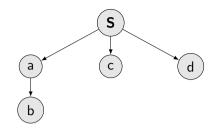


14/33

Which nodes to *explore*? What are the properties of a strategy/algorithm?

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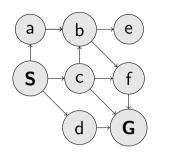


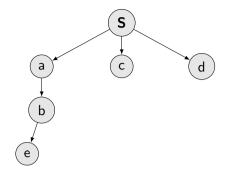


14/33

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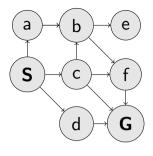
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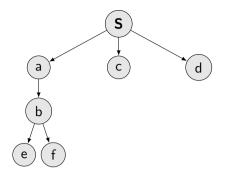




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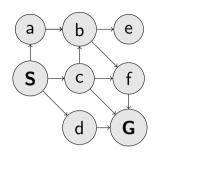


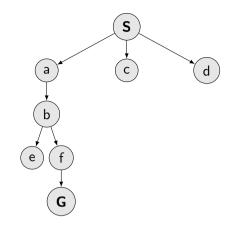


14/33

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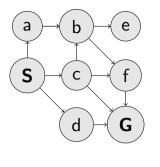


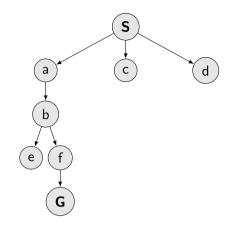
14/33

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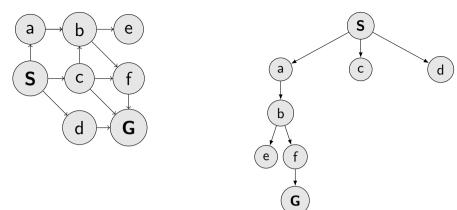


Which nodes to explore?

What are the properties of a strategy/algorithm?

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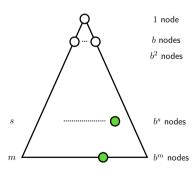
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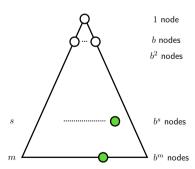
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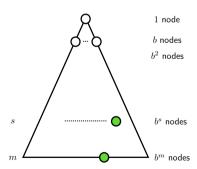
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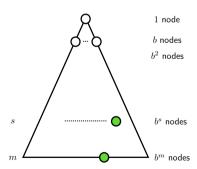
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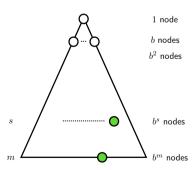
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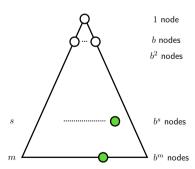
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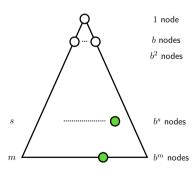
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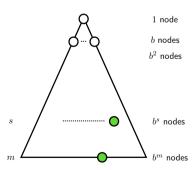
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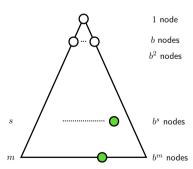
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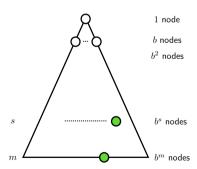
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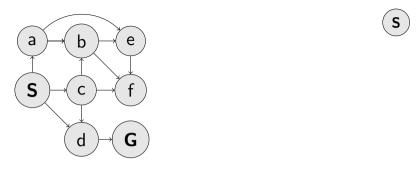


Strategies

How to traverse/build a search tree?

- Notes -

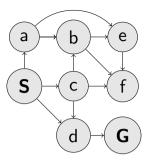
It is perhaps worth to remember that the search tree is built as the algorithm goes. Or better said, the tree is a human friendly representation of the machine run.

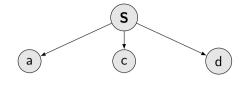


What are the DFS properties (complete, optimal, time, space)?

Notes -

- In animation, we will do the expansion step at once.
- What is the *frontier* set of nodes, waiting to be expanded?
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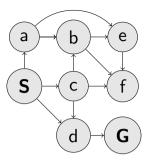


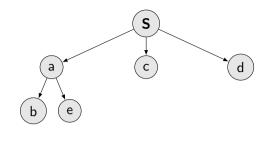


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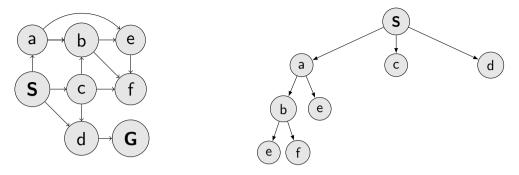




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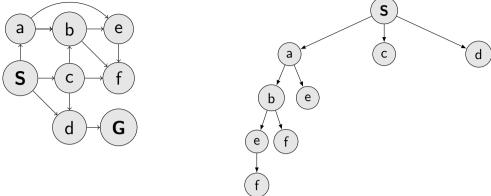
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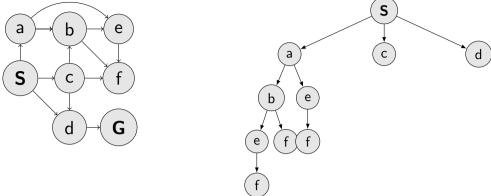
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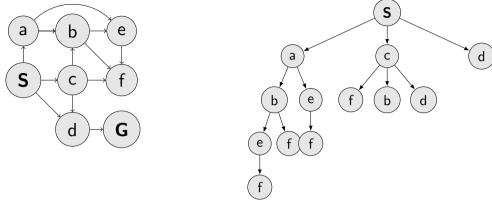
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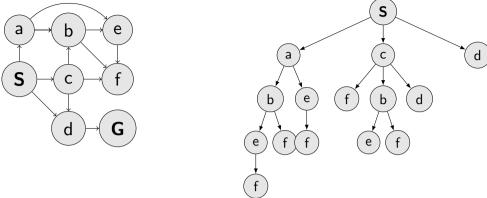
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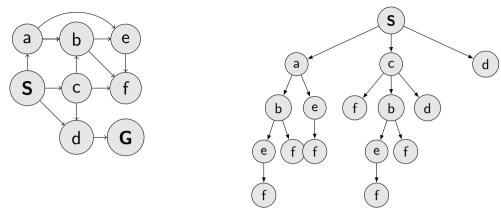
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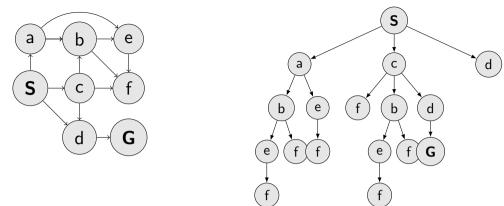
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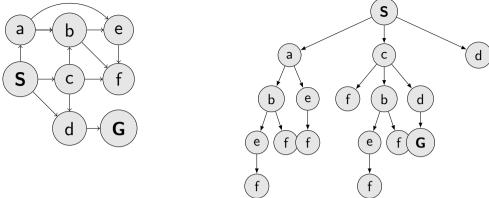
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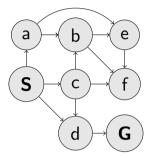
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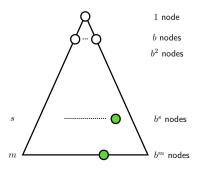
18 / 33

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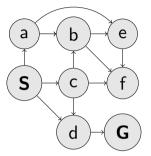


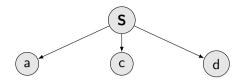
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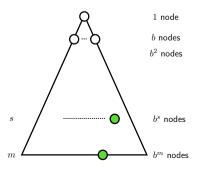
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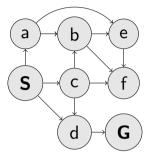


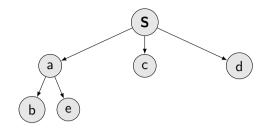
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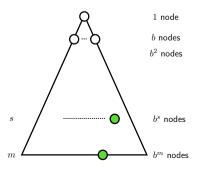
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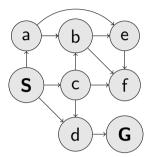


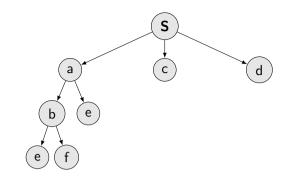
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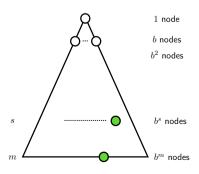
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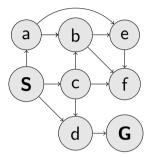
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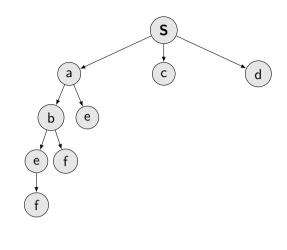
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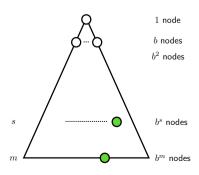
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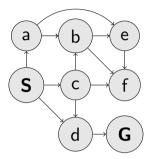
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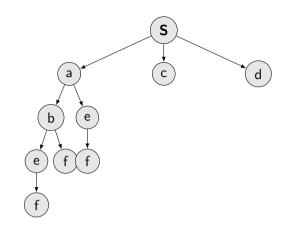
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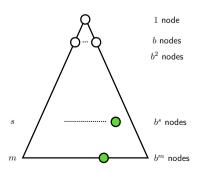
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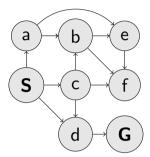
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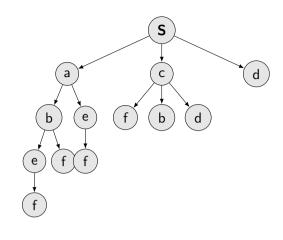
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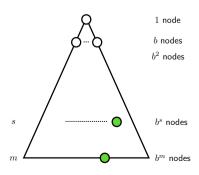
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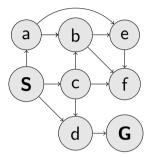
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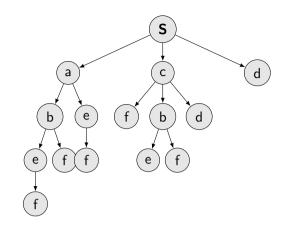
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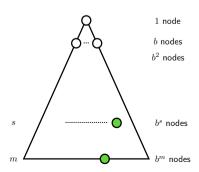
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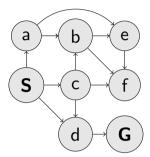


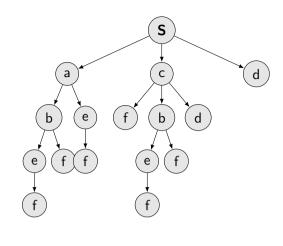
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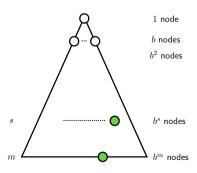
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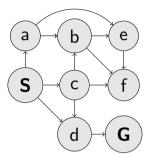
• Space, only the path so far: bm

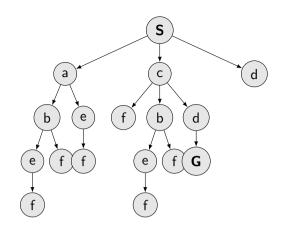
• Completness: m may be ∞ hence, not in general

• Optimality: No! It just takes the first solution found.



- ► Time complexity?
- Space complexity?
- ▶ Complete?
- ► Optimal?





Notes -

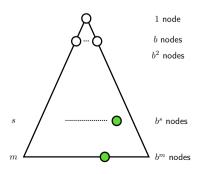
140

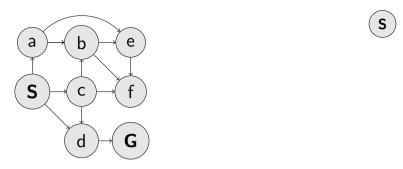
• Space, only the path so far: bm

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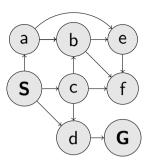
• Completness: m may be ∞ hence, not in general

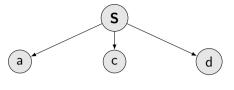
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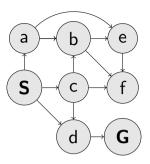


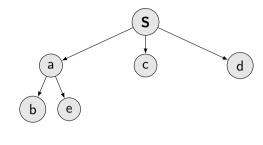
What are the BFS properties?





What are the BFS properties?

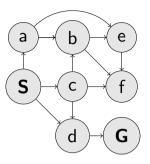


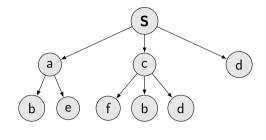


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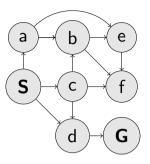
19 / 33

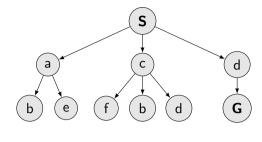
Notes -





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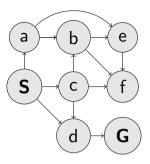


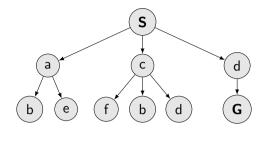


What are the BFS properties?

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Notes -



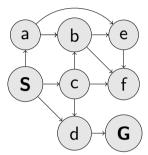


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S

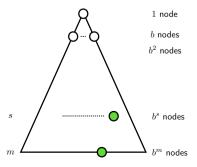
20 / 33

- ► Time complexity?
- ► Space complexity?
- ► Complete?
- ► Optimal?

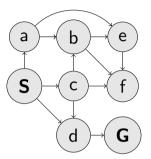


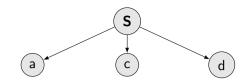
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- Space, all the frontier: b^s
- Completness: Yes!
- Optimality, it does not miss the shallowest solution, hence if all the transition costs are 1: Yes!



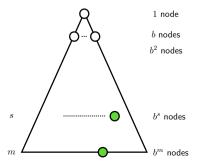
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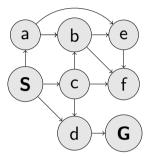


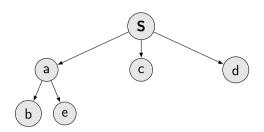
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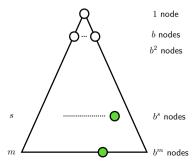
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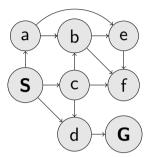


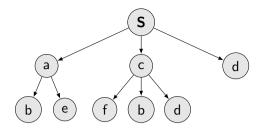
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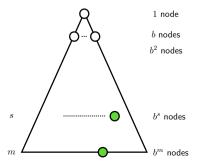
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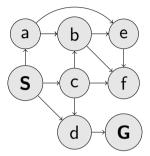


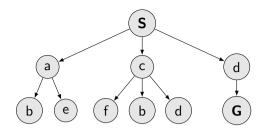
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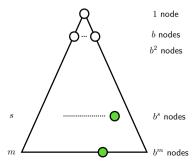
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What are (dis)advantages of the individual strategies?





Notes -

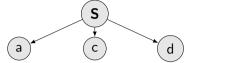
What is the impression from the animation? BFS seems better.

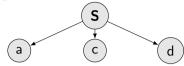
- However, let's not jump to conclusions!
- Draw for yourself a different graph and contruct appropriate trees.
- Not everything is visible from the animations.
- Draw a comparison table.

	Complete	Optimal	Time	Space
DFS	N (Y if no cycles)	N	$O(b^m)$	O(mb)
BFS	Y	Y	$O(b^m)$	$O(b^m)$

- Exponential complexity is scary.
- Practically, space complexity is even more critical. It is not about "waiting longer" but "memory overflow" ...
- This motivates the algorithm modification we look at next.

What are (dis)advantages of the individual strategies?





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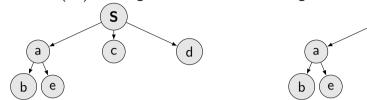
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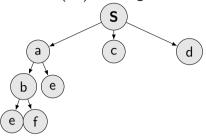
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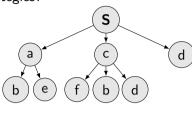
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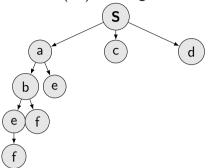
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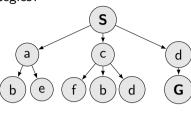
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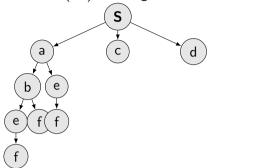
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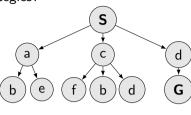
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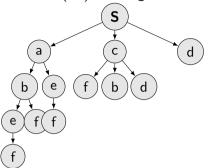
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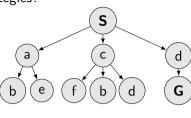
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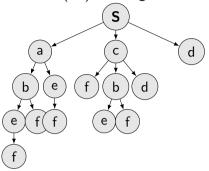
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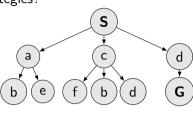
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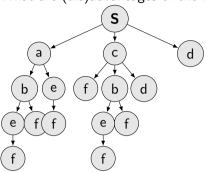
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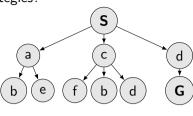
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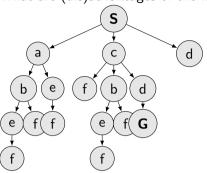
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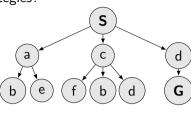
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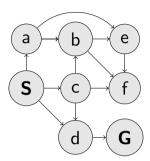
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Do not follow nodes with depth > maxdepth

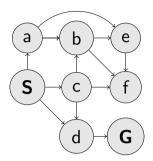


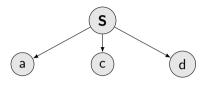


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- Supplying a predetermined max depth nodes at this depth are treated as if they had no successors.
- However, an additional source of algorithm incompleteness. Solution can obviously be deeper than maxdepth unless we know something about the problem. Think about our map of Romania. There are 20 cities. Hence, maxdepth = 19 is a possible choice. Taking a closer look, any city can be reached from any other city in max. 9 steps (state space diameter), giving a more strict and hence better limit.

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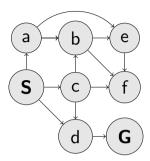


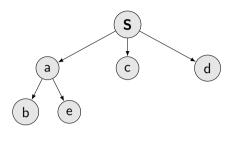


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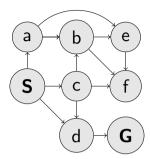


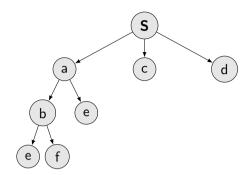


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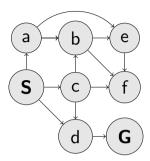


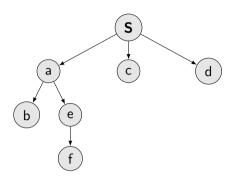


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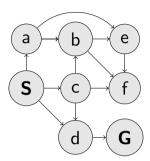


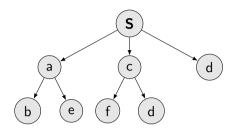


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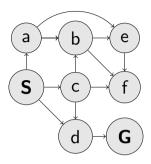


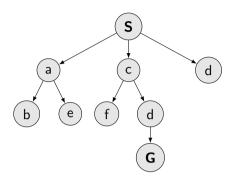
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DFS with limited depth, maxdepth=2

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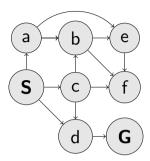


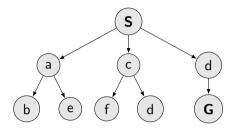
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- However, an additional source of algorithm *incompleteness*. Solution can obviously be deeper than *maxdepth* unless we know something about the problem. Think about our map of Romania. There are 20 cities. Hence, *maxdepth* = 19 is a possible choice. Taking a closer look, any city can be reached from any other city in max. 9 steps (*state space diameter*), giving a more strict and hence better limit.

- ► Start with maxdepth = 1
- ▶ Perform DFS with limited depth. Report success or failure
- ▶ If failure, forget everything, increase maxdepth and repeat DFS

Is it not a terrible waste to forget everything between steps?

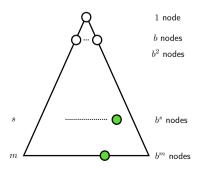
Notes -

Really, how much do we repeat/waste? Compare number of nodes generated ID-DFS vs. BFS:

$$N(ID-DFS) = (s)b + (s-1)b^2 + (s-2)b^3 + \dots + (1)b^s$$

$$N(\mathsf{BFS}) = b + b^2 + b^3 + \dots + b^s$$

Try some calculations for various s and b.



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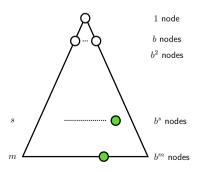
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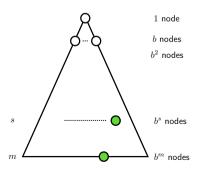
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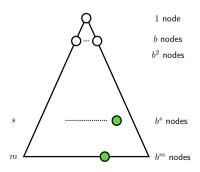
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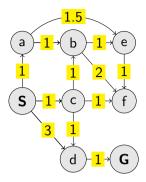
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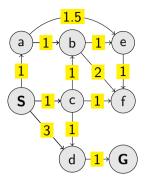


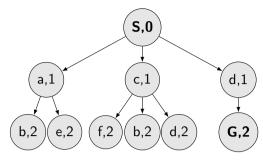




- ▶ In BES_DES_node +depth was the node-value
- ► How was the depth actually computed?
- How to evaluate nodes with path cost?

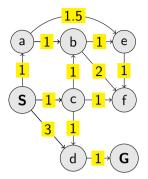
24 / 33

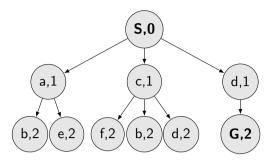




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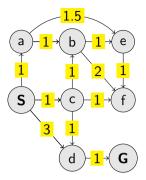
24 / 33

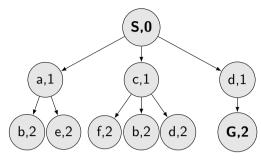




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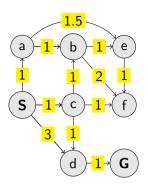
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24 / 33





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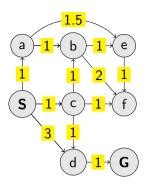
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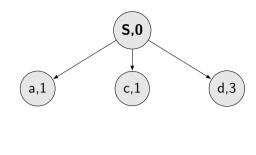
Notes -

Simple extension of BFS. Instead of expanding shallowest node, the node with smallest path cost so far is expanded.

Two differences:

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- Test is added in case a better path is found to a node currently on the frontier.





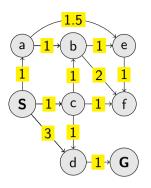
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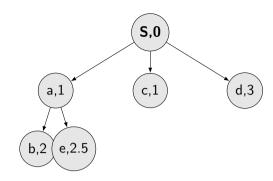
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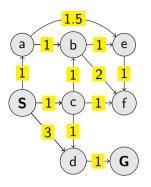
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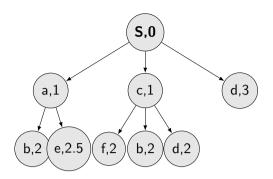
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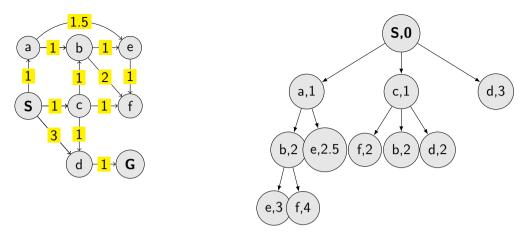
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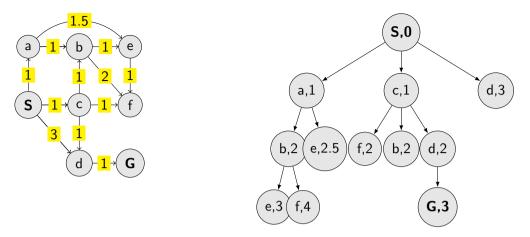
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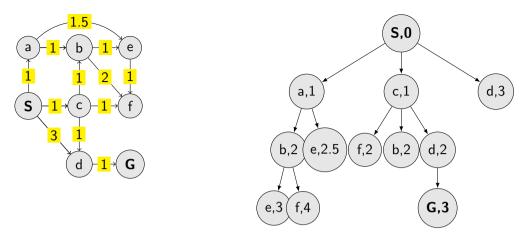
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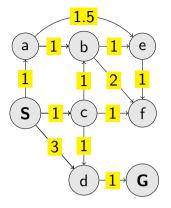
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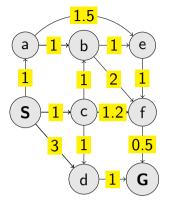
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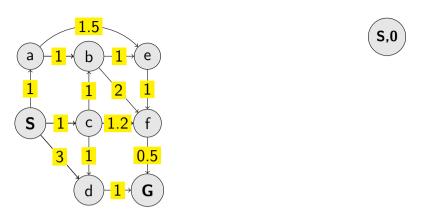
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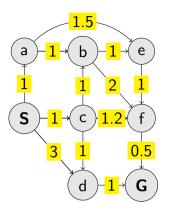
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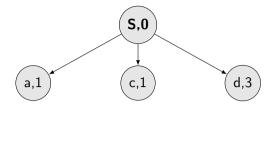




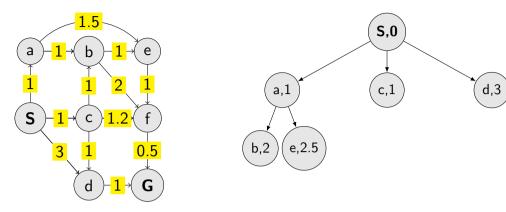
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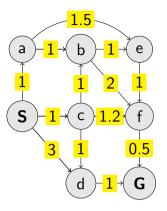


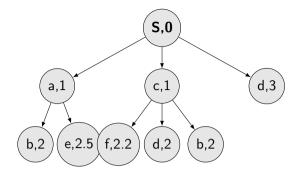


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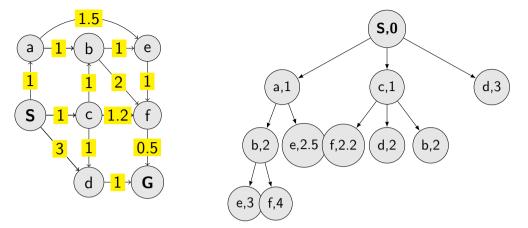


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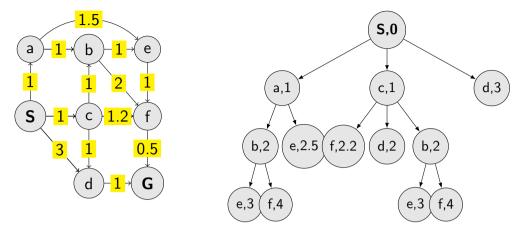




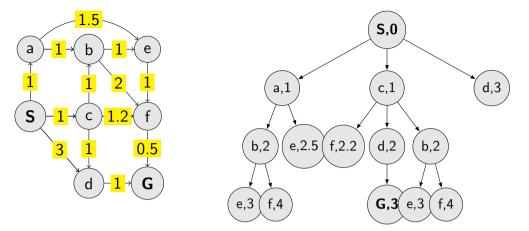
26 / 33



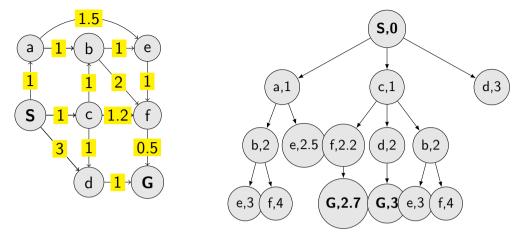
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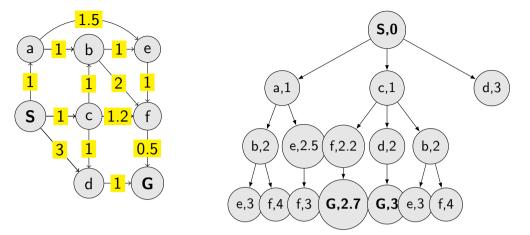
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UCS properties

- ► Time complexity?
- Space complexity?
- ► Complete?
- ► Optimal?

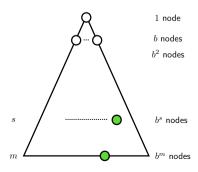
Notes -

Solution cost C^* , transition cost at least ϵ . Effective depth, roughly C^*/ϵ .

• Time: $b^{C^*/\epsilon}$ • Space: $b^{C^*/\epsilon}$

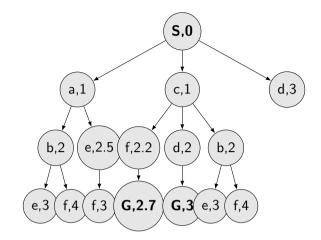
• Completness: Yes!

• Optimality: Yes! Why?



UCS properties

- ► Time complexity?
- Space complexity?
- Complete?
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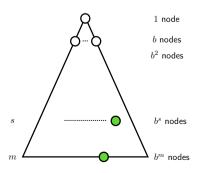
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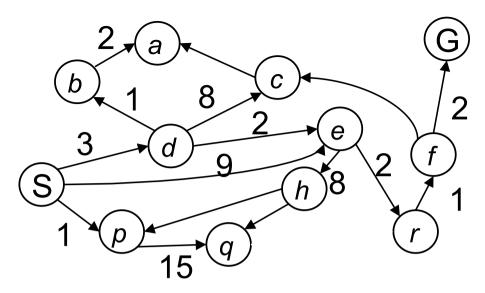
• Space: $b^{C^*/\epsilon}$

• Completness: Yes!

• Optimality: Yes! Why?



Example: Graph with costs



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Notes -

Try it on paper, mark which nodes are in frontier, mark lines of equal cost.



Notes —

Infrastructure for (tree) search algorithms

What should a tree node n now?

- ▶ n.state
- ▶ n.parent
- ▶ n.pathcost

Perhaps we may add something later, if needed . . .

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How to organize nodes?

The Python examples are just suggestions, ...

- ► A dynamically linked structure (list()).
- ► Add a node (list.insert(node)).
- ▶ Take a node and remove from the structure (node=list.pop()).
- ► Check the Python modules heapq¹ and queue² for inspiration.

Notes

https://docs.python.org/3.5/library/heapq.html

²https://docs.python.org/3.5/library/queue.html

What is the solution?

- ▶ We stop when Goal is reached.
- ► How do we contruct the path?

References, further reading

Some figures if from [2]. Chapter 2 in [1] provides a compact/dense intro into search algorithms.

[1] Steven M. LaValle.

Planning Algorithms.

Cambridge, 1st edition, 2006.

Online version available at: http://planning.cs.uiuc.edu.

[2] Stuart Russell and Peter Norvig.

Artificial Intelligence: A Modern Approach.

Prentice Hall, 3rd edition, 2010.

http://aima.cs.berkeley.edu/.